# **PROJECT REPORT**

# MINING ONLINE REVIEWS FOR PRIDICTING SALES PERFORMANCE

# **Department of Computer Science & Engineering**



Under the Supervision of

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Submitted in partial fulfillment of the requirements

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## **BACHELOR OF TECHNOLOGY**

## JAYPEE UNIVERSITY OF INFORMATION TECHNOLOGY

## WAKNAGHAT

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# CERTIFICATE

This is to certify that the work titled "MINING ONLINE REVIEW FOR PRIDICTING SALES PERFORMANCE" submitted by saurabh kumar in partial fulfillment for the award of degree of B.Tech of Jaypee University of Information Technology, Waknaghat has been carried out under my supervision. This work has not been submitted partially or wholly to any other University or Institute for the award of this or any other degree or diploma.

Signature of Supervisor : .....

Name of Supervisor : Dr. Rajni Mohana

Designation :

Date:

# ACKNOWLEDGEMENT

I would like to express our sincere thanks and gratitude to our guide **Dr.Rajni Mohana** for his valuable guidance and suggestions. I am highly indebted to him for providing me an excellent opportunity to learn and present my studies in the form of this project.

I take this opportunity to thank the members of the teaching and non-teaching staff of JUIT for the timely help extended by them.

Lastly thanking our parents, for their morale support and encouragement.

Signature of the student : .....

Name of Student: Saurabh kumar

Date:

# **CHAPTER 1: Introduction**

People's opinion has become one of the extremely important sources for various services in evergrowing popular social networks. In particular, online opinions have turned into a kind of virtual currency for businesses looking to market their products, identify new opportunities, and manage their reputations. Hence, a mechanism that can provide users with condensed descriptions of documents will facilitate the delivery of digital content. This paper explores and designs a mobile system for movie rating and review summarization in which semantic orientation of comments, the limitation of small display capability of cellular devices, and system response time are considered. Practically, when we are not familiar with a specific product, we ask our trusted sources to recommend one.

Today, the popularity of the Internet drives people to search for other people's opinions from the Internet before purchasing a product or seeing a movie. Many websites provide user rating and commenting services, and these reviews could reflect users' opinions about a product. For example, the customer-review section in Amazon.com lists the number of reviews, the percentage for different ratings, and comments from reviewers. When people want to purchase books, CDs, or DVDs, these comments and ratings usually influence their purchasing behaviors. In addition to these websites, a search engine is another important source for people to search for other people's opinions. When a user enters a query into a search engine, the search engine examines its index and provides a listing of best-matching web pages according to its criteria, usually with a short summary containing the document's title and, sometimes, parts of the text.

# **1.1: AIM**

This aim of this project is to develop software that will be able to read the blog or tweet from the web and then it will tell us the reviews based on the sentiments. For this we will approach to design a system that minimizes the ambiguity and redundancy problems. so that people can find out the review of the particular product or the movie etc.

We will try to implement a stemming algorithm, a classifier and a naïve bias to fulfill our project requirement.

# 1.2: EXISTING SYSTEM

A probabilistic latent semantic analysis (plsa) approach has already been presented that calculate the words and then provides the reviews on the basis of calculation. Plsa is a mathematical approach towards the calculation of total words. In plsa it searches for the total words and then calculates for the sentiment of the words and then provide the overall positivity or negativity.

# 1.3: DRAWBACKS OF THE EXISTNG SYSTEM

- Does not convert the word into root word.
- Redundancy is not removed.
- Cannot remove the repeated words.
- Cannot identify the sentiments.

# **CHAPTER 2: TECHNICAL DETAILS**

# **Software Environment**

### Java Technology

Java technology is both a programming language and a platform.

#### The Java Programming Language

The Java programming language is a high-level language that can be characterized by all of the following buzzwords:

- Simple
- Architecture neutral
- Object oriented
- Portable
- Distributed
- High performance
- Interpreted
- Multithreaded
- Robust
- Dynamic
- Secure

With most programming languages, you either compile or interpret a program so that you can run it on your computer. The Java programming language is unusual in that a program is both compiled and interpreted. With the compiler, first you translate a program into an intermediate language called *Java byte codes* —the platform-independent codes interpreted by the interpreter on the Java platform. The interpreter parses and runs each Java byte code instruction on the computer. Compilation happens just once; interpretation occurs each time the program is executed. The following fig. illustrates how this works.

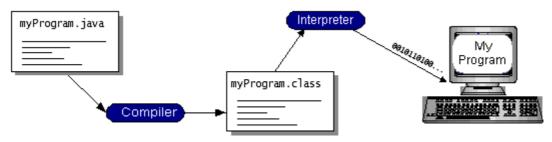


Fig. 1: java environment

You can think of Java byte codes as the machine code instructions for the *Java Virtual Machine* (Java VM). Every Java interpreter, whether it's a development tool or a Web browser that can run applets, is an implementation of the Java VM. Java byte codes help make "write once, run anywhere" possible. You can compile your program into byte codes on any platform that has a Java compiler. The byte codes can then be run on any implementation of the Java VM. That means that as long as a computer has a Java VM, the same program written in the Java programming language can run on Windows 2000, a Solaris workstation, or on an iMac.

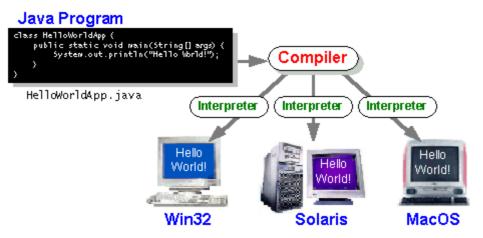


Fig. 2:compiling and processing

## The Java Platform

A *platform* is the hardware or software environment in which a program runs. We've already mentioned some of the most popular platforms like Windows 2000, Linux, Solaris, and MacOS. Most platforms can be described as a combination of the operating system and hardware. The Java platform differs from most other platforms in that it's a software-only platform that runs on top of other hardware-based platforms.

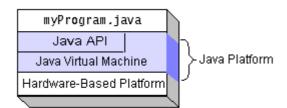
The Java platform has two components:

- The Java Virtual Machine (Java VM)
- The Java Application Programming Interface (Java API)

You've already been introduced to the Java VM. It's the base for the Java platform and is ported onto various hardware-based platforms.

The Java API is a large collection of ready-made software components that provide many useful capabilities, such as graphical user interface (GUI) widgets. The Java API is grouped into libraries of related classes and interfaces; these libraries are known as *packages*. The next section, What Can Java Technology Do? Highlights what functionality some of the packages in the Java API provide.

The following fig. depicts a program that's running on the Java platform. As the fig. shows, the Java API and the virtual machine insulate the program from the hardware.



#### Fig. 3: java extension

Native code is code that after you compile it, the compiled code runs on a specific hardware platform. As a platform-independent environment, the Java platform can be a bit slower than native code. However, smart compilers, well-tuned interpreters, and just-in-time byte code compilers can bring performance close to that of native code without threatening portability.

#### What Can Java Technology Do?

The most common types of programs written in the Java programming language are *applets* and *applications*. If you've surfed the Web, you're probably already familiar with applets. An applet is a program that adheres to certain conventions that allow it to run within a Java-enabled browser.

However, the Java programming language is not just for writing cute, entertaining applets for the Web. The general-purpose, high-level Java programming language is also a powerful software platform. Using the generous API, you can write many types of programs.

An application is a standalone program that runs directly on the Java platform. A special kind of application known as a *server* serves and supports clients on a network. Examples of servers are Web servers, proxy servers, mail servers, and print servers. Another specialized program is a *servlet*. A servlet can almost be thought of as an applet that runs on the server side. Java Servlets are a popular choice for building interactive web applications, replacing the use of CGI scripts. Servlets are similar to applets in that they are runtime extensions of applications. Instead of working in browsers, though, servlets run within Java Web servers, configuring or tailoring the server.

How does the API support all these kinds of programs? It does so with packages of software components that provides a wide range of functionality. Every full implementation of the Java platform gives you the following features:

- The essentials: Objects, strings, threads, numbers, input and output, data structures, system properties, date and time, and so on.
- Applets: The set of conventions used by applets.
- Networking: URLs, TCP (Transmission Control Protocol), UDP (User Data gram Protocol) sockets, and IP (Internet Protocol) addresses.
- **Internationalization**: Help for writing programs that can be localized for users worldwide. Programs can automatically adapt to specific locales and be displayed in the appropriate language.
- **Security**: Both low level and high level, including electronic signatures, public and private key management, access control, and certificates.
- Software components: Known as JavaBeans<sup>TM</sup>, can plug into existing component architectures.
- **Object serialization**: Allows lightweight persistence and communication via Remote Method Invocation (RMI).
- Java Database Connectivity (JDBC<sup>TM</sup>): Provides uniform access to a wide range of relational databases.

The Java platform also has APIs for 2D and 3D graphics, accessibility, servers, collaboration, telephony, speech, animation, and more. The following fig. depicts what is included in the Java 2 SDK.

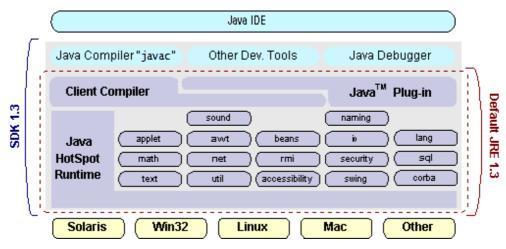


Fig. 4: java database connectivity

### How Will Java Technology Change My Life?

We can't promise you fame, fortune, or even a job if you learn the Java programming language. Still, it is likely to make your programs better and requires less effort than other languages. We believe that Java technology will help you do the following:

- Get started quickly: Although the Java programming language is a powerful object-oriented language, it's easy to learn, especially for programmers already familiar with C or C++.
- Write less code: Comparisons of program metrics (class counts, method counts, and so on) suggest that a program written in the Java programming language can be four times smaller than the same program in C++.
- Write better code: The Java programming language encourages good coding practices, and its garbage collection helps you avoid memory leaks. Its object orientation, its JavaBeans component architecture, and its wide-ranging, easily extendible API let you reuse other people's tested code and introduce fewer bugs.
- **Develop programs more quickly**: Your development time may be as much as twice as fast versus writing the same program in C++. Why? You write fewer lines of code and it is a simpler programming language than C++.
- Avoid platform dependencies with 100% Pure Java: You can keep your program portable by avoiding the use of libraries written in other languages. The 100% Pure Java<sup>TM</sup> Product Certification Program has a repository of historical process manuals, white papers, brochures, and similar materials online.
- Write once, run anywhere: Because 100% Pure Java programs are compiled into machine-independent byte codes, they run consistently on any Java platform.
- **Distribute software more easily**: You can upgrade applets easily from a central server. Applets take advantage of the feature of allowing new classes to be loaded "on the fly," without recompiling the entire program.

### **ODBC**

Microsoft Open Database Connectivity (ODBC) is a standard programming interface for application developers and database systems providers. Before ODBC became a *de facto* standard for Windows programs to interface with database systems, programmers had to use proprietary languages for each database they wanted to connect to. Now, ODBC has made the choice of the database system almost irrelevant from a coding perspective, which is as it should be. Application developers have much more important things to worry about than the syntax that is needed to port their program from one database to another when business needs suddenly change.

Through the ODBC Administrator in Control Panel, you can specify the particular database that is associated with a data source that an ODBC application program is written to use. Think of an ODBC data source as a door with a name on it. Each door will lead you to a particular database. For example, the data source named Sales Fig.s might be a SQL Server database, whereas the Accounts Payable data source could refer to an Access database. The physical database referred to by a data source can reside anywhere on the LAN.

The ODBC system files are not installed on your system by Windows 95. Rather, they are installed when you setup a separate database application, such as SQL Server Client or Visual Basic 4.0. When the ODBC icon is installed in Control Panel, it uses a file called ODBCINST.DLL. It is also possible to administer your ODBC data sources through a stand-alone program called ODBCADM.EXE. There is a 16-bit and a 32-bit version of this program and each maintains a separate list of ODBC data sources.

From a programming perspective, the beauty of ODBC is that the application can be written to use the same set of function calls to interface with any data source, regardless of the database vendor. The source code of the application doesn't change whether it talks to Oracle or SQL Server. We only mention these two as an example. There are ODBC drivers available for several dozen popular database systems. Even Excel spreadsheets and plain text files can be turned into data sources. The operating system uses the Registry information written by ODBC Administrator to determine which low-level ODBC drivers are needed to talk to the data source (such as the interface to Oracle or SQL Server). The loading of the ODBC drivers is transparent to the ODBC application program. In a client/server environment, the ODBC API even handles many of the network issues for the application programmer.

The advantages of this scheme are so numerous that you are probably thinking there must be some catch. The only disadvantage of ODBC is that it isn't as efficient as talking directly to the native database interface. ODBC has had many detractors make the charge that it is too slow. Microsoft has always claimed that the critical factor in performance is the quality of the driver software that is used. In our humble opinion, this is true. The availability of good ODBC drivers has improved a great deal recently. And anyway, the criticism about performance is somewhat analogous to those who said that compilers would never match the speed of pure assembly language. Maybe not, but the compiler (or ODBC) gives you the opportunity to write cleaner programs, which means you finish sooner. Meanwhile, computers get faster every year.

#### **JDBC**

In an effort to set an independent database standard API for Java; Sun Microsystems developed Java Database Connectivity, or JDBC. JDBC offers a generic SQL database access mechanism that provides a consistent interface to a variety of RDBMSs. This consistent interface is achieved through the use of "plug-in" database connectivity modules, or *drivers*. If a database vendor wishes to have JDBC support, he or she must provide the driver for each platform that the database and Java run on.

To gain a wider acceptance of JDBC, Sun based JDBC's framework on ODBC. As you discovered earlier in this chapter, ODBC has widespread support on a variety of platforms. Basing JDBC on ODBC will allow vendors to bring JDBC drivers to market much faster than developing a completely new connectivity solution.

JDBC was announced in March of 1996. It was released for a 90 day public review that ended June 8, 1996. Because of user input, the final JDBC v1.0 specification was released soon after.

The remainder of this section will cover enough information about JDBC for you to know what it is about and how to use it effectively. This is by no means a complete overview of JDBC. That would fill an entire book.

#### **JDBC Goals**

Few software packages are designed without goals in mind. JDBC is one that, because of its many goals, drove the development of the API. These goals, in conjunction with early reviewer feedback, have finalized the JDBC class library into a solid framework for building database applications in Java.

The goals that were set for JDBC are important. They will give you some insight as to why certain classes and functionalities behave the way they do. The eight design goals for JDBC are as follows:

#### 1. SQL Level API

The designers felt that their main goal was to define a SQL interface for Java. Although not the lowest database interface level possible, it is at a low enough level for higher-level tools and APIs to be created. Conversely, it is at a high enough level for application programmers to use it confidently. Attaining this goal allows for future tool vendors to "generate" JDBC code and to hide many of JDBC's complexities from the end user.

#### 2. SQL Conformance

SQL syntax varies as you move from database vendor to database vendor. In an effort to support a wide variety of vendors, JDBC will allow any query statement to be passed through it to the underlying database driver. This allows the connectivity module to handle non-standard functionality in a manner that is suitable for its users.

3. *JDBC must be implemental on top of common database interfaces* The JDBC SQL API must "sit" on top of other common SQL level APIs. This goal allows JDBC to use existing ODBC level drivers by the use of a software interface. This interface would translate JDBC calls to ODBC and vice versa.

#### 4. Provide a Java interface that is consistent with the rest of the Java system

Because of Java's acceptance in the user community thus far, the designers feel that they should not stray from the current design of the core Java system.

#### 5. Keep it simple

This goal probably appears in all software design goal listings. JDBC is no exception. Sun felt that the design of JDBC should be very simple, allowing for only one method of completing a task per mechanism. Allowing duplicate functionality only serves to confuse the users of the API.

#### 6. Use strong, static typing wherever possible

Strong typing allows for more error checking to be done at compile time; also, less error appear at runtime.

#### 7. Keep the common cases simple

Because more often than not, the usual SQL calls used by the programmer are simple SELECT'S, INSERT'S, DELETE'S and UPDATE'S, these queries should be simple to perform with JDBC. However, more complex SQL statements should also be possible.

Finally we decided to proceed the implementation using Java Networking.

And for dynamically updating the cache table we go for MS Access database.

Java has two things: a programming language and a platform.

Java is a high-level programming language that is all of the following

Simple	Architecture-neutral
Object-oriented	Portable
Distributed	High-performance
Interpreted	multithreaded
Robust	Dynamic
Secure	

Java is also unusual in that each Java program is both compiled and interpreted. With a compile you translate a Java program into an intermediate language called Java byte codes the platform-independent code instruction is passed and run on the computer.

Compilation happens just once; interpretation occurs each time the program is executed. The fig. illustrates how this works.

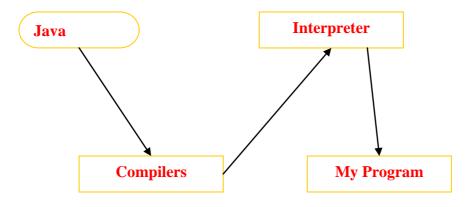


Fig. 5: java program

You can think of Java byte codes as the machine code instructions for the Java Virtual Machine (Java VM). Every Java interpreter, whether it's a Java development tool or a Web browser that can run Java applets, is an implementation of the Java VM. The Java VM can also be implemented in hardware.

Java byte codes help make "write once, run anywhere" possible. You can compile your Java program into byte codes on my platform that has a Java compiler. The byte codes can then be run any implementation of the Java VM. For example, the same Java program can run Windows NT, Solaris, and Macintosh.

#### Map Visualizations

Charts showing values that relate to geographical areas. Some examples include: (a) population density in each state of the United States, (b) income per capita for each country in Europe, (c) life expectancy in each country of the world. The tasks in this project include:

Sourcing freely redistributable vector outlines for the countries of the world, states/provinces in particular countries (USA in particular, but also other areas);

Creating an appropriate dataset interface (plus default implementation), a rendered, and integrating this with the existing XYPlot class in JFreeChart;

Testing, documenting, testing some more, documenting some more.

#### Time Series Chart Interactivity

Implement a new (to JFreeChart) feature for interactive time series charts --- to display a separate control that shows a small version of ALL the time series data, with a sliding "view" rectangle that allows you to select the subset of the time series data to display in the main chart.

#### Dashboards

There is currently a lot of interest in dashboard displays. Create a flexible dashboard mechanism that supports a subset of JFreeChart chart types (dials, pies, thermometers, bars, and lines/time series) that can be delivered easily via both Java Web Start and an applet.

#### **Property Editors**

The property editor mechanism in JFreeChart only handles a small subset of the properties that can be set for charts. Extend (or re implement) this mechanism to provide greater end-user control over the appearance of the charts.

# **System Configuration**

# H/W System Configuration:

Processor - Penti	um –III
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Speed	- 1.1 Ghz	
RAM	- 256 MB(min)	
Hard Disk	- 20 GB	
Floppy Drive	- 1.44 MB	
Key Board	- Standard Windows Keyboar	ď
Mouse	- Two or Three Button Mous	e
Monitor	- SVGA	

# S/W System Configuration:

Operating System	: Windows95/98/2000/XP
Application Server	: Tomcat5.0/6.X
Front End	: HTML, Java, Jsp
Scripts	: JavaScript.
Server side Script	: Java Server Pages.
Database	: Mysql
Database Connectivity	: JDBC.C

# **CHAPTER 3: ARCHITECTURE**

The design of the proposed system discusses the various requirements that will make up the application. Installation on Windows requires WinPcap software which can be downloaded from winPcap website. Jpcap is a set of Java classes which provide an interface and system for network packet capture, it is required for packet capture in Java and built upon Libpcap which is a packet capture library in C language. Java Runtime Environment (JRE) 5.0 or higher will also be required to run this Java application. More space may be required to store the captured packets since the required space on hard disk for installation is less than 1MB.

By conducting the requirements analysis we listed out the requirements that are useful to restate the problem definition

# Data Flow Diagram / Use Case Diagram / Flow Diagram

The DFD is also called as bubble chart. It is a simple graphical formalism that can be used to represent a system in terms of the input data to the system, various processing carried out on these data, and the output data is generated by the system.

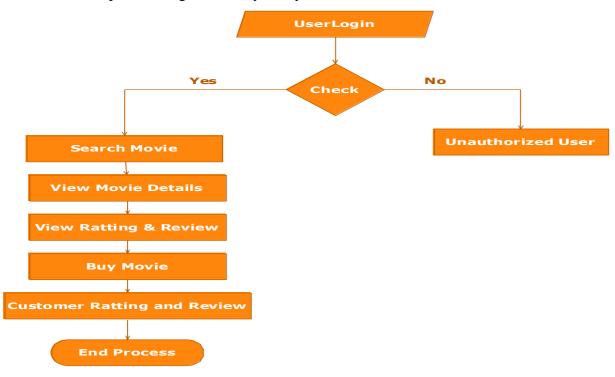


Fig. 6: data flow diagram(user)

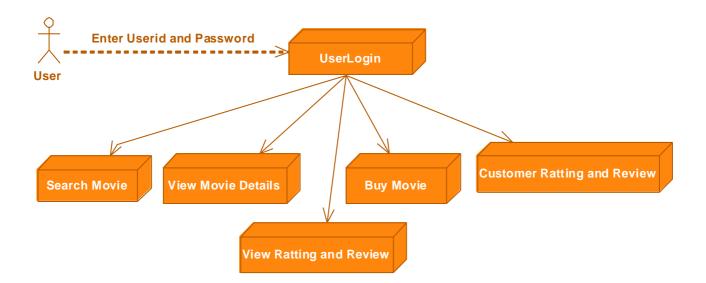


Fig. 7: component diagram

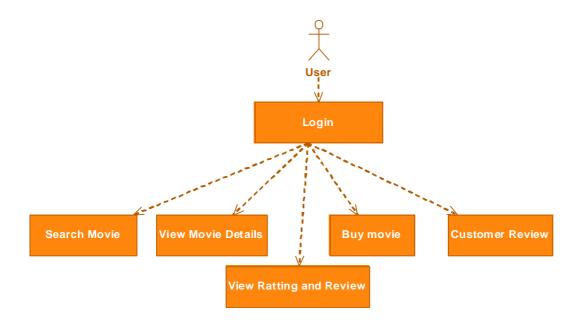


Fig. 8: use case diagram

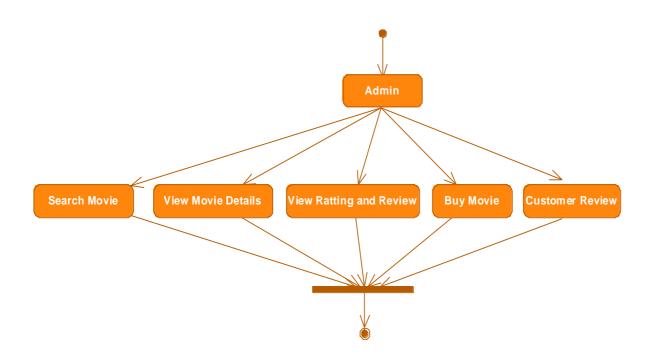


Fig. 9: activity diagram

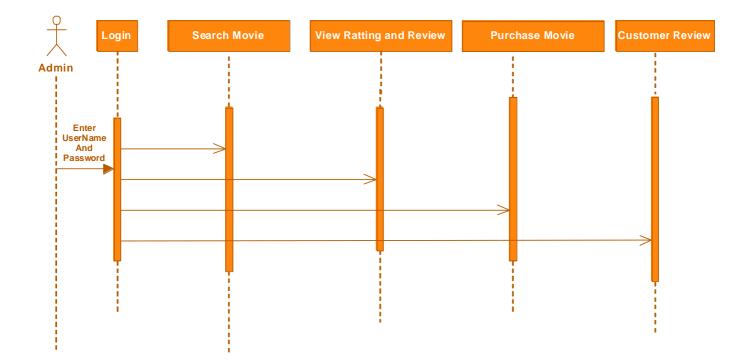


Fig. 10: sequence diagram

# **CHAPTER 4: SYSTEM DESIGN**

### **4.1:** MODULES:

- 1. Sentiment Analysis
- 2. Feature-Based Summarization
- 3. Dataset
- 4. Opinion-Word Identification

# **Modules Description**

#### 1. Sentiment Analysis

Since a document is composed of sentences and a sentence is composed of terms, it is reasonable to determine the semantic orientation of the text from terms. As a result, the sentiment analysis research started from the determination of the semantic orientation of the terms. Hatzivassiloglou and McKeown employed textual conjunctions such as "fair and legitimate" or "simplistic but well-received" to separate similarly connoted and oppositely connoted words. Esuli and Sebastiani proposed to determine the orientation of subjective terms based on the quantitative analysis of the glosses of such terms, i.e., the textual definitions that are given in online dictionaries. The process is based on the assumption that terms with similar orientation tend to have "similar" glosses (i.e., textual definitions). Thus, synonyms and antonyms could be used to define a relation of orientation. Esuli and Sebastiani described SENTIWORDNET, which is a lexical resource in which each WordNet synset is associated with three numerical scores, i.e., Obj(s), Pos(s), and Neg(s), thus describing how objective, positive, and negative the terms contained in the synset.

#### 2. Feature-Based Summarization

In product-review summarization, people are interested in the reasons why this product is worth buying rather than the principal meaning of the comment. Thus, feature-based summarization is used in movie-review summarization. The feature-based summarization will focus on the product features on which the customers have expressed their opinions. In addition to product features, the summarization should include opinion information about the product; therefore, product features and opinion words are both important in feature-based summarization. As a result, product features and opinion-word

identification are essential in feature-based summarization.

#### 3. Product-Feature Identification

We propose an LSA-based product-feature-identification algorithm and system can obtain a semantically related feature set for each seed. We compared three product-feature-identification approaches, i.e., ratting about product feature, price and delivery.

#### 4. Opinion-Word Identification

In addition to feature identification, opinion words about the product features are important as well. Hu and Liu extracted the opinion words by retrieving the nearby adjective of product features. In addition to language sentence-structure characteristic, Zhuang *et al.* used the dependency grammar graph to find out some relations between feature words and the corresponding opinion words in training data. They both rely on language sentence structure to extract opinion words; therefore, these approaches will be applicable to those language sentences having such a characteristic.

## **4.2: SYSTEM STUDY**

### **4.2.1: FEASIBILITY STUDY**

The feasibility of the project is analyzed in this phase and business proposal is put forth with a very general plan for the project and some cost estimates. During system analysis the feasibility study of the proposed system is to be carried out. This is to ensure that the proposed system is not a burden to the company. For feasibility analysis, some understanding of the major requirements for the system is essential.

Three key considerations involved in the feasibility analysis are

- ♦ ECONOMICAL FEASIBILITY
- ♦ TECHNICAL FEASIBILITY
- ♦ SOCIAL FEASIBILITY

### ECONOMICAL FEASIBILITY

This study is carried out to check the economic impact that the system will have on the organization. The amount of fund that the company can pour into the research and development of the system is limited. The expenditures must be justified. Thus the developed system as well within the budget and this was achieved because most of the technologies used are freely available. Only the customized products had to be purchased.

#### **TECHNICAL FEASIBILITY**

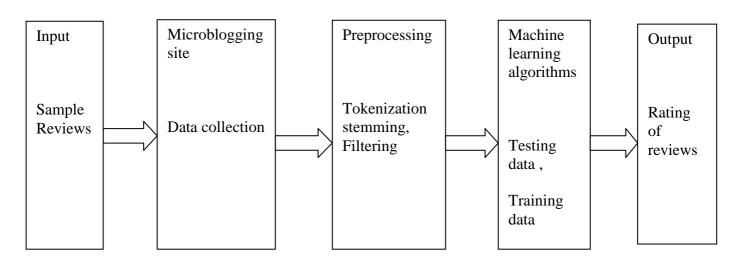
This study is carried out to check the technical feasibility, that is, the technical requirements of the system. Any system developed must not have a high demand on the available technical resources. This will lead to high demands on the available technical resources. This will lead to high demands being placed on the client. The developed system must have a modest requirement, as only minimal or null changes are required for implementing this system.

#### SOCIAL FEASIBILITY

The aspect of study is to check the level of acceptance of the system by the user. This includes the process of training the user to use the system efficiently. The user must not feel threatened by the system, instead must accept it as a necessity. The level of acceptance by the users solely depends on the methods that are employed to educate the user about the system and to make him familiar with it. His level of confidence must be raised so that he is also able to make some constructive criticism, which is welcomed, as he is the final user of the system.

# **CHAPTER 5: METHODOLOGY**

Methodologies used :





#### 5.1: Data collection:

The first work in this thesis is to collect huge amount of data collection of data although is not a simple task as it may seem to anyone at first thought. For this we have to make certain assumption and decision. Basically we have to collect two different data sets: test data, training data.

**Test data Collection:** One of the objective of this thesis is to analyze the sentiment of messages posted in reaction to movie reviews, so we have to collect tweets about movies only. However, it is not a simple task.

A corpus of around 1000 tweets about movie reviews was sampled from the test data and we manually examined and annotated that data as positive or negative because the data which we collected was unlabelled ,so we formed two classes as pos, neg.

Out of thousand post that have been annotated 31 posts here needed context to understand and determine their sentiments. thus 97% of our test data did not require context to determine their sentiments. Many a times in a time came when it was difficult to determine the sentiment, as they seemed both positive and negative.

#### 4. Training data:

We have chosen only subjective data for training ,no objective was collected here.

Subjective data involves positive and /or negative sentiment only. if does not consists of any neutral tweets.

A total of 5000 subjective twitter posts were collected for corpus. One this data was collected, it was separated into tweets that contain only negative emoticons tweets that contain only positive emoticons and tweets that contain both emoticons.

Class	Training data	Test data
Negative	2000	400
Positive	3000	600

#### Table 1: data collection

#### **5.3:Supervised approaches**

This is where most of time and effort was spent in our thesis under this section, different supervised machine learning approaches were used and explained. To be able to experiment with different machine learning algorithms and to able the identification of factors that affect our results ,we used a three step process.

The first step is the preprocessing of the training data. The second is feature extraction and data representation in a particular format.(semantic vsm here).

The third is the classifier training and testing phase with different machine learning algorithms . this approaches helped in experimenting with different possibilities by varying the preprocessing operations, feature extraction and then using the machine learning algorithms .each of this will be explained in the following subsections.

#### **5.3.1:Pre-processing:**

#### **Cleaning of the data:**

Since the corpus we are taking consists of syntactical features that may not be useful for machine learning the data needs to be cleaned for further use

Tokenization: in this process, we have broken the stream of text up into words, phrases, symbols and then called them as token the frequency of a particular token in the whole data set is also counted. With that we have also taken out total no. of unique words in our data set.

Stemming: in this process we have reduced the inflicted words to their stem or root word. For this we have taken in to account that stem need not to be identically same to the morphological root of the word ; it is sufficient that a related words could map to the same stem is not itself a valid root Stop word removal : after tokenization we are having a big list of words occurring in data sets and mostly not all of them are useful for learning task. It is imperative to reduce the size of feature space as far as possible. So stop word removal step will be done prior too tokenization to remove all occurrence of these useless words like 'a', 'an', 'the', 'is', etc.

Filtering : during filtering process various sub tasks were done .

url :url were removed ,in order to remove the feature size during training

emoticons : all emoticons were replaced by pos and neg classes

**username and hash tags :** we replaced them with *<*un> and *<*ht> respectively.

#### **Removal of repeated comments**

**Lower casing :** all characters were lower cased to ensure that all tokens map to the corresponding feature irrespective of casing .

Class	Data	After duplicate	Duplicates
		removal	
Positive	3000	2960	40
Negative	2000	1990	10
Total	5000	4950	50

**5.3.2: Removing duplicates:** duplicates were removed from all training data as they just add Weight and nothing else .only exact duplicates were removed by us .the result of the removal of duplicates are given in table as can be seen from the table there were duplicates present in each class . the no. of duplicates is more in positive tweets then there are in negative tweets ,because people have tendency to re tweet something they are very happy with the total percentage of duplicates is 10% only.

**5.3.3:Feature extraction and instance representation :** the most important question in machine learning is how to represent data both the training and test data must be represented in some or the other way in order for a machine learning algorithm to learn and thus build a model out of it .some of the ways in which the data can be represented could be feature based or bag of words representation .by features based ,it is meant that some attributes which we think that they capture the pattern of the data are selected first and the entire data set must be represented n terms of them before it is fed to any machine learning algorithm.

Different features like n-gram presence or n-gram frequency ,pos tags ,syntactic features or semantic features can be used here in this representation .

This paper explores and designs a system for movie rating and review summarization in which semantic orientation of comments. System response time are considered. Practically, when we are not familiar with a specific product, we ask our trusted sources to recommend one. Today, the popularity of the Internet drives people to search for other people's opinions from the Internet before purchasing a product or seeing a movie. Many websites provide user rating and commenting services, and these reviews could reflect users' opinions about a product. For example, the customer-review section in Amazon.com lists the number of reviews, the percentage for different ratings, and comments from reviewers.

When people want to purchase books, CDs, or DVDs, these comments and ratings usually influence their purchasing behaviors. In addition to these websites, a search engine is another important source for people to search for other people's opinions. When a user enters a query into a search engine, the search engine examines its index and provides a listing of best-matching web pages according to its criteria, usually with a short summary containing the document's title and, sometimes, parts of the text. In this paper, we collected movie reviews from Internet Blogs that do not consist of any rating information. Sentiment analysis is performed to determine the semantic orientation of the reviews and movie-rating score is based on the sentiment-analysis result.

## 5.3.4: Stemming:

A stemming algorithm is a computational procedure which reduces all words with the same root (or, if pre-fixes are left untouched, the same stem) to a common form, usually by stripping each word of its derivational and inflectional suffixes. Researchers in many areas of computational linguistics and information retrieval find this a desirable step, but for varying reasons. In automated morphological analysis, the root of a word may be of less immediate interest than its suffixes, which can be used as clues to grammatical structure. (See, e.g., Earland Resnikoff and Dolby .This field has also been reported on by S. Silver and M. Lott, Machine Translation Project, University of California, Berkeley [personal communication].) At the other extreme, what suffixes are found may be subsidiary to the problem of removing them consistently enough to obtain sets of exactly matching stems. Word-frequency counts using stems, for stylistic (as described by S. Y. Sedelow [per- sonal communication]) or mathematical analysis of a body of language, often require matched stems. (So does stemming as part of an information-retrieval sys- tem, the specific application which motivated this pa- per. ) But certain linguistic problems are common to any "stem-oriented" stemming algorithm, no matter what its ultimate use. The brief description below of the framework within which Project Interrex is planning to use a stemming algorithm should be viewed as but one possible application for research on the morphological structure of English and other languages. Similarly, a variety of applications are considered in evaluating the theoretical and practical attributes of several previous algorithms. As a major part of its information transfer experiments, Project Intrex is developing an integrated re- trieval system in which a library user, through a remote computer terminal, can first obtain extensive information from a central digital store about documents that are available on a specific subject, and then obtain the full text of the documents. A prototype retrieval system is being assembled in order to permit experimentation with its various components. The experimental system will use a specially compiled augmented library catalogue containing information on approximately 10,000 documents in the field of materials science and engineering, including not only author, title, and other basic data about each document but also an abstract, bibliography, and a list of subject terms indicating the content of the document. Each subject term is a phrase of one or more English words. A stemming algorithm will be used to maximize the usefulness of the subject terms.

In many cases, the information which is semantically significant to the user of the system is contained in the stems of the lexical words in the subject terms, and suffixes and function words merely enable this information to be expressed in a grammatical form. The form of the words which the user inputs will often not correspond to that of the original words in the catalogue. To Permit the words in the user's query to match the words in the catalogue entry's subject terms, both query and subject terms can be stripped of the suffixes that prevent their matching. For example, computational and computing might both be stemmed to compute.

In constructing the software needed for this particular application of stemming (or any other), we encounter questions which are answerable only in terms of the over-all system. For instance, what should constitute a "word" to be stemmed? In the case of Intrex, what suffixes should the algorithm search for that are specifically 22 oriented toward terms in materials science and engineering? These are questions of less general interest than the linguistic problems of extracting a stem from any one word in a non-specialized vocabulary (for an example of lists of affixes taken from terms in specific technical fields, see Dyson). The development of an efficient algorithm should logically precede investigation of these questions, and they will not be discussed further here. The approach to stemming taken here involves a two- phase stemming system. The first phase, the stemming algorithm proper, retrieves the stem of a word by removing its longest possible ending which matches one on a list stored in the computer. The second phase handles "spelling exceptions," mostly instances in which the "same" stem varies slightly in spelling

according to what suffixes originally followed it. For example, absorption will be output from phase one as absorpt, absorbing as absorb. The problem of the spelling exceptions, which in the above example involves matching absorpt and absorb, is discussed thoroughly in Section V of this paper. One particular solution to the problem, termed recoding, has been implemented in the present phase two. We also plan to use the present basic algorithm as a foundation in testing out other feasible solutions.1 This plan is appropriate because spelling-ex- ception rules can, and probably should, be formulated independently of the stemming algorithm proper.

II. Stemming, Form, and Meaning By its computational nature, a stemming algorithm has inherent limitations. The routine handles individual words: it has no access to information about their gram- matical and semantic relations with one another. In fact, it is based on the assumption of close agreement of meaning between words with the same root. This assumption, while workable in most cases, in English represents an approximation at best. It is a better or worse approximation depending on the intended use of the stems, the semantic vagaries of individual roots, and the strength of the algorithm (how radically it transforms words). A stemming algorithm strong enough to group together all words with the same root may be unsuitable for, say, word-frequency counting. For such applications one would not wish a pair like neutron-.neutral- izer to coincide, and one would prefer to work with a very limited list of suffixes. Where stems are used as a means of associating related items of information, as they are in an automated library catalogue, and where the catalogue can be interrogated in an on-line mode, it seems best to use a strong algorithm, that is, one that will combine more words into the same group rather than fewer, thus pro- viding more document references rather than fewer. I am indebted to Richard S. Marcus and Peter Kugel for valuable discussion of this specific problem and of this report as a whole. After a word in the library user's query has been stemmed and a matching stem and associated list oriented toward terms in materials science and engineering? These are questions of less general interest than the linguistic problems of extracting a stem from any one word in a non-specialized vocabulary (for an example of lists of affixes taken from terms in specific technical fields, see Dyson ). The development of an efficient algorithm should logically precede investigation of these questions, and they will not be discussed further here.

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Where stems are used as a means of associating re- lated items of information, as they are in an automated library catalogue, and where the catalogue can be in- terrogated in an on-line mode, it seems best to use a strong algorithm, that is, one that will combine more words into the same group rather than fewer, thus providing more document references rather than fewer. After a word in the library user's query has been stemmed and a matching stem and associated list of ull-word forms has been found in the catalogue and presented to the user, he may decide to discard some of these forms in order to inhibit searching for those full- word forms which are unrelated to his subject.

Occasionally, the output of a stemming routine may be not only ambiguous but also "not English." This happens when a suffix is identical to the end of some root. For instance, -ate is a noun suffix in directorate, but simply part of a verbal root in create and appreciate. In English, situations of this type limit the use of suf- fixes as clues to parts of speech. Sometimes grammatical information is required for stemming, not provided by it. However, the generation of such non linguistic stems as cre - and appreci- is not a serious problem; if the pur- pose of stemming is only to allow related words to match, then the stems yielded by a stemming algorithm need not coincide with those found by a linguist. The exact form of the stem is not critical if it is the same no matter what suffixes have been removed following it, and if "mistaken" stemming does not generate an am-

biguity. Similarly, the ending that must be removed in order to achieve a consistent algorithm is determined in relation to the stemming system as a whole. The ending may or may not be exactly equivalent to some entity in English morphology, and it may be acceptable to have the computer program remove it when a linguist would not, with no detriment to the ultimate results.

#### **III.** Types of Stemming Algorithms

Two main principles are used in the construction of a stemming algorithm: iteration and longestmatch. An algorithm based solely on one of these methods often has drawbacks which can be offset by employing some combination of the two principles.

Iteration is usually based on the fact that suffixes are attached to stems in a "certain order, that is, there exist order-classes of suffixes (see, e.g., Lejnieks [4]). Each order-class may or may not be represented in any given word. The last order-class—the class that occurs at the very end of a word—contains inflectional suffixes such as -s, -es, and-ed. Previous order-classes are derivational.(As pointed out by J. L. Dolby [personal communication], there are several cases known in which a derivational suffix (-ness) follows an inflectional one (-ed or -ing). This occurs with certain nominalized adjectives derived from verbs by use of one of these two inflectional endings, for example, relatedness, disinterestedness, willingness.) An example of the lowest order-class in a word may be what is technically part of the root (see the -ate example above), but for the purposes of computation it is considered part of the ending. An iterative stemming algorithm is simply a recursive procedure, as its name implies, which removes strings in each order-class one at a time, starting at the end of a word and working toward its beginning. No more than one match is allowed within a single order-class, by definition. One must decide how many order-classes there should be, which endings should occur in each, and whether or not the members of each class should be internally ordered for scanning.

The longest-match principle states that within any given class of endings, if more than one ending provides a match, the one which is longest should be removed. This principle is implemented by scanning the endings in any class in order of decreasing length. For example, if -ion is removed when there is also a match on -ation, provision would have to be made to remove -at, that is, for another order-class. To avoid this extra order-class, -ation should precede -ion on the list. An algorithm based strictly on the longest-match principle uses only one order-class. All possible combinations of affixes are compiled and then ordered on length. If a match is not found on longer endings, shorter ones are scanned. The obvious disadvantage to this method is that it requires generating all possible combinations of affixes. A second disadvantage is the amount of stor- age space the endings require.

The first disadvantage may also be present to a large degree when one is setting up an iterative algorithm with as many order-classes as possible. To set up the order-classes, one must examine a great many endings. Furthermore, it is not always obvious to which class a given string should belong for maximum efficiency. It is also entirely possible that the occurrence of members of some classes is context dependent (see below). In short, while an iterative algorithm requires a shorter list of endings, it introduces a number of complications into the preparation of the list and programming of the routine.

Some idea of the breadth of these complications is gained through consideration of another basic attribute of a stemming algorithm: it is context free or context sensitive. Since "context" is used here to mean any attribute of the remaining stem, "context free" implies

no qualitative or quantitative restrictions on the removal of endings. In a context-free algorithm, the first ending in any class which achieves a match is accepted. But

here should presumably be at least some quantitative restriction, in the sense that the remaining stem must not be of length zero. An example of this extreme case is the matching of -ability to ability as well as to com- putability. In fact, any useful stem usually consists of at least two letters, and often three or four constitute a necessary minimum. The restriction on stem length varies with the ending; how it varies can again only be determined in relation to the total system. The algorithm developed by Professor John W. Tukey of Prince- ton University (personal communication) associates a lower limit with each ending. Some of his limits are quite high (e.g., seven letters). I have been less conservative and have proposed a minimum stem length of two; certain endings have an additional restriction in that their minimum stem length is three, four, or five letters.

The kind of qualitative contextual restrictions that should be imposed is a somewhat open question. In order to get the best results, certain endings should not be removed in the presence of certain letters in the resultant stem, usually those letters that immediately pre- cede the ending. The more desirable form of context- sensitive rule is a general one that can be applied to a number of endings, but such rules are few. One example is "do not remove an ending that begins with -en-,following -e." Violation of this rule would change seen to se-, a potentially ambiguous stem (cf. sea minus -a, seize minus- ize, etc.). But a number of rules must be created for individual endings in order to avoid certain special cases peculiar to those endings. One can go to great lengths in this direction, with increasingly small returns. I have preferred to start by treating a number of the more obvious exceptions in the hope that the percentage of words not accounted for will be small enough to preclude the need to add many additional rules. An iterative stemming algorithm, that is, one that contains more than one order-class of endings, is presumably no less complicated by

context-sensitive rules than a one-class algorithm, and is probably more so; exceptions associated with the members of each class may depend on a rather complicated context. For ex- ample, suppose there is a rule (in a non-iterative algorithm) stating that minimal stem length is five before -ionate. The endings -ion and -ate occur separately, also, with different restrictions. In an iterative routine, -ion and -ate would only occur as separate endings, in different order-classes; and -ion would be restricted by the rule that its preceding context must be of length five if -ate was found during the preceding iteration. In other words, the endings that are removed may influence the lower-order endings that can be removed subsequently. The implications for simplicity in programming are self-evident. In a pure longest-match algorithm, the only context that need be considered is the prospective stem itself.

Since computer-storage space for endings was not an immediate problem, it was decided to test a non-iterative stemming algorithm based on a one-class list of endings. That is, the intuitively inefficient procedure of listing both singular and plural forms, and so on, has been followed in order to minimize the number of con- text-sensitive rules necessary. Compilation of the actual list of endings used is discussed in the next section; the algorithm is outlined in Section VI.

The author is aware of three previous major attempts to construct stemming algorithms. Tukey has proposed a context-sensitive, partially iterative stemming algorithm whose endings are divided into four order-classes. The first (highest-order) class contains only terminal s which, however, is not removed after i, s, or u. The second class is recursive, the third is non-recursive and ordered on length. The fourth class consists of remaining terminal consonants. The last three classes also have a few members each with simple context restrictions, and all classes have limits on minimum stem length. (The basic structure of this "tail-cropping" algorithm is not affected by its multilingual orientation, though the endings used would obviously differ from those found in a procedure for English only.) One of the more interesting things about the Tukey system is its structural complexity. One class uses the longest-match principle only, while another is iterative (and thus not a proper order-class). Presumably the object of this heterogeneous structure is to avoid the repetitiveness of a one-class ending list in the most concise way possible. However, as stated earlier, there is a compromise between conciseness of rules and simplicity of programming. By contrast, the algorithm developed at Harvard University by Michael Lesk, under the direction of Professor Gerard Salton [10], is based on an iterated search for a longest-match ending. After no more matches can be found, terminal i, a, and e are removed, and then possibly terminal consonants. There are apparently no con- textual restrictions of any kind. (A brief description of the algorithm, including a useful list of 194 endings, was transmitted to us via personal communication. A sample of these suffixes, and further information about the algorithm, have more recently appeared in Salton.)

A third algorithm has been developed by James L. Dolby of R and D Consultants, Los Altos, California (personal communication). This algorithm works in three stages, the first of which involves a set of context- dependent transformations. Most of the cropping is done in the second stage, a context-free, longest-match, re- cursive procedure which removes endings in any order but is subject to the restriction of a two-syllable mini- mum stem length. In the final stage there is a context dependent dropping of inflectional forms. The endings used were derived by algorithm from word lists on the basis of orthographic context, and are "minimal" segments of one to four letters in length.

#### IV. Compilation of a List of Endings

A one-class list of endings (concatenations of suffixes) was compiled in the following way: A preliminary list was based on endings found in a small portion of the augmented catalogue being developed by Project Interrex and on endings in the list used at Harvard. The preliminary list was evaluated by applying the endings on this list to a portion of the output from Tukey's tail- cropping routine, levels 1-3, and volumes 5-7 of the Normal and Reverse English Word List [8] (volumes 5 7 contain unbroken words sorted alphabetically when written from right to left). Since each of these lists is organized according to ends of words, it was possible to see whether the removal of a given ending would result in (1) two different stems matching, or (2) a stem not matching another stem which it should match. Either of these conditions, unless it was caused by a spelling exception or caused improper matching in only a few rare cases, necessitated the addition of new endings, the disposing of old ones, or the addition of context- may occur, but do not always occur. The second assumption is that these changes involve no more than two letters at the end of a stem--this is merely an empirical result which has not yet been contradicted. It has also been observed that the sequences of letters that cause difficulty are often common to more than one class of exceptions. In recoding, this means that some rules can cover more than one type of exception, although it is not usually the case.

The crucial difference between recoding and partial matching is this: a recoding procedure is part of the stemming algorithm while a partial-matching procedure is not. Partial matching operates on the output from the stemming routine at the point where the stems derived from catalogue terms are being searched for matches to the user's stemmed query. All partial matches, within certain limits, are retrieved rather than just all perfect matches; discrepancies are resolved after retrieval, not in the previous stemming procedure. This has the advantages of reducing stemming to the onestep process of removing an ending and of eliminating the context specifications sometimes needed in recoding. The dis- advantages, which are not so obvious, can be discussed only after a more complete description of a partial matching procedure is given. Such a procedure starts with an unmodified stem S1

—again, absorption is a good example. The first step is to search the list of stemmed catalogue terms for all those which begin with S1 minus its last two letters: in this

case, all stems of any length beginning with absor, which we call S2. Of course, special provisions will have to be made for cases in which S1 is only two or three letters long. Among those stems returned will be absorpt and absorb. Absorbefaci, the stem of absorbefacient, may also be found. This last item will be eliminated, probably for the better, by the next step of the procedure, which discards all stems more than two characters longer than S1 (here, more than nine letters long). We then have collected all stems which match absorpt within two letters in either direction. Given any one of these, Sj, a final match is allowed between Sj and S1 if and only if either Sj = S1 or the following conditions are satisfied:

1. The stems Sj and S1 must match at least up to two letters before the end of the longer of them.

2. If Sj and S1 are the same length and differ by one letter, this letter plus a blank must occur on a closed list (see Appendix D) for each stem.

3. If Sj and S1 are the same length and differ by two letters, each sequence of two letters must occur on the list.

4. If Sj and S1 differ in length by one, the last two letters of the longer, and the last of the shorter plus a blank, must occur on the list.

5. If Sj and S1 differ in length by two, the last two letters of the longer must occur on the list. The above rules amount essentially to examining the last two letters of stems that match up to that point; if the stems are different lengths, all "missing letters" in the shorter are represented by blanks. The "closed list" needed for this routine is given in Appendix D.

It may appear that an unacceptable number of "wrong" matches would result from this procedure, since there are no restrictions on which pairs of items on the list may be used to produce a match. There are two defenses against this view:

First, such a closed list does exist. Many partial matches will *not* be allowed. Of those that are allowed erroneously, many would have been produced also by a recoding procedure, for much the same reasons.

Second, we can make a probabilistic argument. Most of the stems used will probably be fairly

long—long enough so that there are unlikely to be many matches within two letters. Any  $S_j$  found by searching with  $S_2$  stands a good chance of being related to  $S_2$ , and thus to  $S_1$ .

In short, while a partial-matching procedure may produce *no fewer wrong* matches than recoding, it will probably produce *more right* ones. It is inherently more flexible than recoding rules; all classes of exceptions do not have to be specified beforehand. Part of this flexi-bility results from allowing  $S_1$  and  $S_j$  to differ in length by two letters in *either* direction. Yet this condition also provides a built-in barrier against certain types of wrong matches, as the following example illustrates:

*Convex* is recoded to convic by the rule  $ex \rightarrow ic$ ; *convict*, the stem of *conviction*, is recoded to *convict by* the rule  $ct \rightarrow c$ . This erroneous match is *not* allowed in partial matching, since although condition (4) is satisfied, condition (1) is not.

Partial matching is a kind of controlled recoding; the recoding takes place *only* if a partial, but not complete, match is found. The original stem is still preserved, however, providing a constant check for violation of condition.

Using partial matching as a substitute for recoding does have one major disadvantage for a system using disk storage, as Interrex does, and it is a potentially serious one. In some cases, the timeconsuming retrieval from the disk of a great number of partial matches, those beginning with  $S_2$ , will be necessary. These cases are most likely to occur with very short stems. The question is whether in such instances  $S_2$  can be lengthened (made closer to  $S_1$ ) enough to avoid this problem and still retrieve all acceptable matches. Empirical data are needed to answer this question, as well as to determine whether the number of short stems used is great enough to warrant concern. Any timing, programming, or other complications which partial matching introduces must be small enough to be balanced out by other advantages it may offer.

#### VI. The Two-Phase Stemming Routine and Its Results

Several progressively more advanced versions of the Interrex stemming routine have been coded in AED (a compiler language developed at the Electronic Systems Laboratory) [7, pp. *367-85]* and run on sample batches of words, using the MIT 7094 CTSS system. The flow chart in Fig. 1 shows the most important features of the stemming and recoding parts of the program.

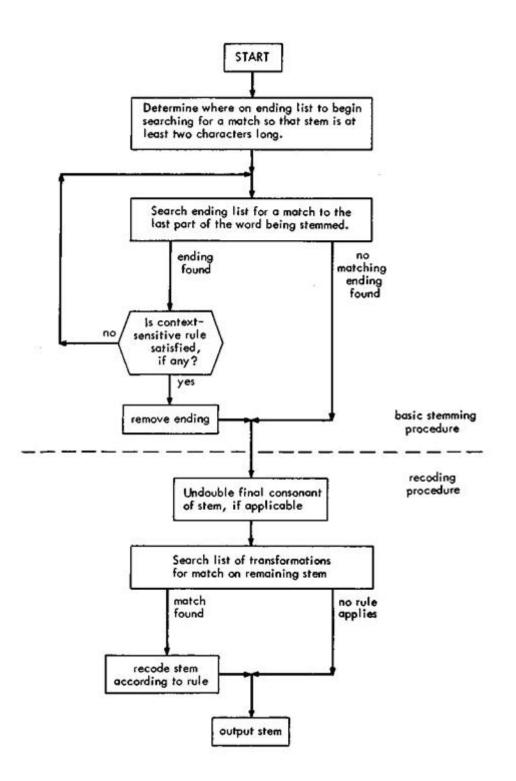


Fig. 12: porter's stemming

While a full evaluation of this stemming system within the Project Intrex environment will not be possible until the augmented catalogue data base is completed, output so far indicates that the procedures used are workable and will yield very good results with only minor changes. These changes involve the list of endings and occasionally the recoding rules; the types of operations performed remain the same.

To give some idea of the alterations that are needed to make the system highly effective, I shall discuss several of the changes that have been made in the pro-gram. Fig. 2 shows the result of stemming several groups of related words. An obvious problem was that "magnet" and "magnesium" had the same recoded stem. This problem was easy to fix by changing recoding rule 32 from  $et \rightarrow es$  to  $et \rightarrow es$  except following *n*.

An additional recoding rule took care of the discrepancy between *meter*  $\rightarrow$  *meter* and *metric*  $\rightarrow$  *metr:metr* 

 $\rightarrow$  meter. All other changes involved the stemming procedure: *-ium*, *-ite*, and *-itic* were added to the list of endings, with the stipulation that *-ite* be removed only in certain rather limited cases and *-itic* only after t or ll; the rule governing *-al-* endings was changed so that they are not removed after met-; l was added to the list of stem-final consonants to be undoubted; and the context in which the removal of *-on* is allowable was broadened to include single t. The results after these changes are shown in Fig. 3. It is expected that several more such evaluations of a random group-sample will catch most similar difficulties still left in the program, although it is likely that minor revisions will be required as long as the vocabulary of the data base continues to increase.

## **Porter's stemming :**

The ANSI C version that heads the table below is exactly equivalent to the original BCPL version. The BCPL version did, however, differ in three minor points from the published algorithm and these are clearly marked in the downloadable ANSI C version. They are discussed further below.

This ANSI C version may be regarded as definitive, in that it now acts as a better definition of the algorithm than the original published paper.

Over the years, I have received many encoding from other workers, and they are also presented below. I have a reasonable confidence that all these versions are correctly encoded.

language	author	affiliation	received	notes
<u>ANSI C</u>	me			
ANSI C thread safe	Me			
<u>java</u>	Me			
<u>Perl</u>	Me			
<u>Perl</u>	Daniel van Balen		Oct 1999	slightly faster?
<u>python</u>	<u>Vivake</u> <u>Gupta</u>		Jan 2001	
<u>Csharp</u>	<u>André</u> <u>Hazelwood</u>	The Official Web Guide	Sep 2001	
<u>Csharp</u> .NE T compliant	<u>Leif</u> <u>Azzopardi</u>	Univerity of Paisley, Scotland	Nov 2002	

<u>Common</u> <u>Lisp</u>	<u>Steven M.</u> <u>Haflich</u>	Franz Inc	Mar 2002	
Ruby	Ray Pereda	www.raypereda.com	Jan 2003	<u>github link</u>
<u>Visual</u> <u>Basic</u> VB6	<u>Navonil</u> <u>Mustafee</u>	Brunel University	Apr 2003	
<u>Delphi</u>	Jo Rabin		Apr 2004	
<u>Javascript</u>	'Andargor'	<u>www.andargor.com</u>	Jul 2004	substantial revisions by <u>Christophe</u> <u>I</u> <u>McKenzie</u> (See 'Links' below.)
Visual Basic VB7; .NET compliant	<u>Christos</u> <u>Attikos</u>	University of Piraeus, Greece	Jan 2005	
<u>php</u>	Richard Heyes	www.phpguru.org	Feb 2005	
<u>Prolog</u>	<u>Philip</u> <u>Brooks</u>	University of Georgia	Oct 2005	
<u>Haskell</u>	<u>Dmitry</u> <u>Antonyuk</u>		Nov 2005	
<u>T-SQL</u>	Keith Lubell	www.atelierdevitraux.com	May 2006	
matlab	Juan Carlos	California Pacific Medical	Sep 2006	

	<u>Lopez</u>	Center Research Institute		
<u>Tcl</u>	<u>Aris</u> <u>Theodorakos</u>	NCSR Demokritos	Nov 2006	
<u>D</u>	<u>Daniel</u> <u>Truemper</u>	Humboldt-Universitaet zu Berlin	May 2007	
<u>erlang</u> (1) erlang (2)	<u>Alden Dima</u>	National Institute of Standards and Technology, Gaithersburg, MD USA	Sep 2007	
<u>REBOL</u>	Dale K Brearcliffe		Apr 2009	
<u>Scala</u>	<u>Ken</u> Faulkner		May 2009	
<u>sas</u>	Antoine St- Pierre	Business Researchers, Inc	Apr 2010	
<u>plugin vim</u> <u>script</u>	<u>Mitchell</u> Bowden		May 2010	<u>github link</u>
node.js	Jed Parsons	jedparsons.com	May 2011	<u>github link</u>
Google Go	<u>Alex</u> Gonopolskiy		Oct 2011	<u>github link</u>
<u>awk</u>	<u>Gregory</u> <u>Grefenstette</u>	3ds.com/exalead	Jul 2012	
<u>clojure</u>	<u>Yushi Wang</u>		Mar 2013	bitbucket link

<u>Rust</u>	<u>Do Nhat</u> <u>Minh</u>	<u>Nanyang</u> <u>University</u>	Technological	Aug 2013	github link
<u>vala</u>	Serge Hulne			Sep 2013	

All these encodings of the algorithm can be used free of charge for any purpose. Questions about the algorithms should be directed to their authors, and not to Martin Porter (except when he is the author).

To test the programs out, here is a <u>sample vocabulary</u> (0.19 megabytes), and the corresponding <u>output</u>.

Email any comments, suggestions, queries

## Points of difference from the published algorithm

There is an extra rule in Step 2, $(m>0) \log i \rightarrow \log$ So archaeology is	equated	with archaeological etc.
The Step 2 rule $(m>0) abli \rightarrow able$		
is replaced by		
$(m>0) bli \rightarrow ble$		
So <i>possibly</i> is	equated	with <i>possible</i> etc.

The algorithm leaves alone strings of length 1 or 2. In any case a string of length 1 will be unchanged if passed through the algorithm, but strings of length 2 might lose a final s, so as goes to a and is to i.

These differences may have been present in the program from which the published algorithm derived. But at such a great distance from the original publication it is now difficult to say.

It must be emphasised that these differences are very small indeed compared to the variations that have been observed in other encodings of the algorithm.

Status

The Porter stemmer should be regarded as 'frozen', that is, strictly defined, and not amenable to slightly further modification. As a stemmer, it is inferior to the Snowball English or Porter2stemmer, which derives from it, and which is subjected to occasional improvements. For practical work, therefore, the new Snowball stemmer is recommended. The Porter stemmer is appropriate to IR research work involving stemming where the experiments need to be exactly repeatable.

## **Common errors**

Historically, the following shortcomings have been found in other encodings of the stemming algorithm.

The algorithm clearly explains that when a set of rules of the type

 $(condition)S1 \rightarrow S2$ 

are presented together, only one rule is applied, the one with the longest matching suffix SI for the given word. This is true whether the rule succeeds or fails (i.e. whether or not S2 replaces S1). Despite this, the rules are sometimes simply applied in turn until either one of them succeeds or the list runs out.

This leads to small errors in various places, for example in the Step 4 rules

	$(m>1)ement \rightarrow$	-	-	-
	$(m>1)ment \rightarrow$			
	$(m>1)ent \rightarrow$			
to		remove		final ement, ment and ent.

Properly, *argument* stems to *argument*. The longest matching suffix is *-ment*. Then stem *argu*-has measure *m* equal to 1 and so *-ment* will not be removed. End of Step 4. But if the three rules are applied in turn, then for suffix *-ent* the stem *argum*-has measure *m* equal to 2, and *-ent* gets removed.

The more delicate rules are liable to misinterpretation. (Perhaps greater care was required in explaining them.) So

((m>1) and (\*s or \*t))ion is taken to mean

(m>1)(s or t)ion

The former means that taking off *-ion* leaves a stem with measure greater than 1 ending *-s* or *-t*; the latter means that taking off *-sion* or *-tion* leaves a stem of measure greater than 1. A similar confusion tends to arise in interpreting rule 5b, to reduce final double L to single L. Occasionally cruder errors have been seen It is interesting that although the published paper explains how to do the tests on the strings SI by a program switch on the last or last but one letter, many encodings fail to use this technique, making them much slower than they need be.

Table 3: Word clusters and stems

Cluster no.	Word	Stem	Cluster no.	Word	Stem
	create	creat			traba
				trabajan	
	creates	creat			traba
			2	trabajar	
	creating	creat	(Spanish)		traba
1				trabajado	
(English)	created	creat			traba
				trabajador	
	creation	creat		dificil	dific
			3		
	creative	creat	(Portuguese)	dificilmente	dific

Now, let us compare our stemmer empirically with an established linguistic stemmer, say Porter's stemmer. We select English as our testing language as outputs of Porter's algorithm can be easily extracted at Snowball project. We take sample of 100 random English words and apply Porter's Snowball framework and our procedure over them. The results are reported in Table 3.2. Identical stems generated by either algorithm are shown in bold. Four POS, Noun, Adjective, Verb and Adverb were taken into account. We select *Newspaper* search area in COCA.

To compare the stemming results of Snowball (Porter) stemmer and our N-gram stemmer, we employed direct evaluation method. Like the comparison made by Chris D. Paise , using length truncation as a baseline, we employed *Levenshtein distance* as a base measure. The distance is the number of deletions, insertions, or substitutions required to transform the source string into the target string. The Levenshtein distance therefore, represents by how many units a word has been stripped by a stemmer. For our sample of the 100 words, these distances are shown in Table 3.2,  $LD_{Snow}$  and  $LD_{N-gram}$  for Snowball stems and our N-gram stems respectively.

## Table 4: Random Words and their stems

					LDN-
S. No.	Word	Snowball stem	N-gram ster	m LDSnov	v gram
1	Damage	D	Description	1	1
	Parsons	Parson	Parson	1	-
2	Dilution	Dilut	Dilut	3	3
3	Agreement	Agreement	Agree	0	4
4	Passion	Passion	Passion	0	0
5	Cutter	Cutter	Cutter	0	0
6	Refill	Refill	Refil	0	1
7	Museums	Museum	Museum	1	1
8	Stallion	Stallion	Stallion	0	0
9	Braid	Braid	Brai	0	1
10	Fleet	Fleet	Flee	0	1
11	Midwife	Midwif	Midwife	1	0
12	Haze	Haze	Haze	0	0
13	Prophet	Prophet	Prophe	0	1
14	Vacancy	Vacanc	Vacanc	1	1
15	Ninety	Nineti	Ninet	1	1
16	Menace	Menac	Menac	1	1
17	Laceration	Lacer	Lacerat	5	3
18	Training	Train	Train	3	3
19	Librarian	Librarian	Librar	0	3
20	Ordaining	Ordain	Ordain	3	3
21	Mainstream	Mainstream	Mainst	0	4
22	Bloodier	Bloodier	Blood	0	3
23	Abject	Abject	Abject	0	0
24	Substitute	Substitut	Substitut	1	1
25	Overseas	Oversea	Overseas	1	0

27 Despotic Despot Despoti 2	1
28 Predicted Predict Predic 2	3
29 Lyric Lyric 0	0
30 Eleventh Eleven 0	2
31 Admonishing Admonish Admoni 3	5
<b>32</b> Quiet <b>Quiet Quiet</b> 0	0
33MacabreMacabrMaca1	3
34UnloadedUnloadUnload2	2
35 Quizzical Quizzic Quizzi 2	3
36AltoAlto0	0
37 Undeveloped Undevelop Undevelo 2	3
38CasualCasual0	0
39RecognizedRecognRecogni4	3

40	Devastating	Devast	Devastat	5	3
41	Abused	- Abus	Abuse	2	1
42	Abusing	Abus	Abusing	3	0
43	Admired	Admir	Admire	2	1
44	Admiring	Admir	Admir	3	3
45	Believed	Believ	Belie	2	3
46	Believing	Believ	Belie	3	4
47	Borrowed	Borrow	Borrow	2	2
48	Borrowing	Borrow	Borrow	3	3
49	Carried	Carri	Carrie	2	1
50	Carrying	Carri	Carry	3	3
51	Consulted	Consult	Consult	2	2
52	Consulting	Consult	Consult	3	3
53	Deceived	Deceiv	Deceive	2	1
54	Deceiving	Deceiv	Deceiv	3	3
55	Decorated	Décor	Décor	4	4

56	Decorating	Décor	Décor	5	5
57	Employed	Employ	Employ	2	2
58	Employing	Employ	Employ	3	3
59	Explained	Explain	Explain	2	2
60	Explaining	Explain	Explain	3	3
61	Finished	Finish	Finish	2	2
62	Finishing	Finish	Finish	3	3
63	Forbidden	Forbidden	Forbid	0	3
64	Forbidding	Forbid	Forbid	4	4
65	Gathered	Gather	Gather	2	2
66	Gathering	Gather	Gather	3	3
67	Improved	Improv	Improv	2	2
68	Improving	Improv	Improv	3	3
69	Laughed	Laugh	Laugh	2	2
70	Laughing	Laugh	Laugh	3	3
71	Listened	Listen	Listen	2	2
72	Listening	Listen	Listen	3	3
73	Mended	Mend	Mende	2	1
74	Mending	Mend	Mending	3	0
75	Nipped	Nip	Nippe	3	1
76	Nipping	Nip	Nipp	4	3
77	Plucked	Pluck	Pluck	2	2
78	Plucking	Pluck	Pluck	3	3
79	Preached	Preach	Preach	2	2
80	Preaching	Preach	Preach	3	3
81	Enormously	Enorm	Enorm	5	5
82	Monthly	Month	Month	2	2
83	Solemnly	Solemn	Solemn	2	2
84	Abnormally	Abnorm	Abnormal	4	2
85	Diligently	Dilig	Diligen	5	3
86	Jubilantly	Jubil	Jubil	5	5

87	Frightfully	Fright	Fright	5	5
88	Swiftly	Swift	Swift	2	2
89	Miserable	Miser	Miser	4	4
90	Thankfully	Thank	Thankful	5	2
91	Blissfully	Bliss	Blissful	5	2
92	Reluctantly	Reluct	Reluctan	5	3
93	Viciously	Vicious	Vicious	2	2
94	Wonderfully	Wonder	Wonder	5	5
95	Hopelessly	Hopeless	Hopeless	2	2
96	Briskly	Brisk	Briskly	2	0
97	Delightfully	Delight	Delight	5	5
98	Anxiously	Anxious	Anxious	2	2
99	Obnoxiously	Obnoxi	Obnoxious	5	2
100	Inwardly	Inward	Inward	2	2

We hypothesize that if our stemmer is as efficient as that of Porter's, then the distributions of Levenshtein distances between a word and its stripped stem tend to be same. For a given word, we have two treatments in hand, Porter's algorithm and our N-gram procedure. Thus for a given word, we have a pair of LDs. For comparison, we set the null hypothesis  $H_0$  that no differences exist between the two stemmers.

Because the two sets of measures can be considered as two measures associated to the same sample, we decide to employ a statistical test for paired samples. In particular, a nonparametric statistical test, the *Wilcoxon signed-rank test* is employed by us, since there was no evidence about the distribution of these distances. The Wilcoxon test is based on two paired series  $(x_i, y_i)$  of *N* observed values, one series for each out of the observed variables *X*, *Y* to be compared. The absolute differences  $d_i = abs (x_i \cdot y_i)$  and  $sign (d_i)$  of differences are to be computed next. Then the absolute values are ranked, discarding zeros. If  $N_r$  is the reduced sample size, the final test statistic is  $W = |\sum_{i=1}^{N} r[sign(d_i) \cdot rank(|d_i|)]|$ , which is approximated by a Normal

variable for large N.

# 4. Experimental Results:

On conducting the Wilcoxon test over our sample data, we get the *p* value of 0.54. As *p*>0.05, fail to reject the null hypothesis  $H_0$ . Thus, we have strong evidence that our N-gram Stemmer is not inferior to Porter's Stemmer.

## Naïve bayes classifier :

A **Bayes** classifier is a simple probabilistic <u>classifier</u> based on applying <u>Bayes'</u> <u>theorem</u> (from <u>Bayesian statistics</u>) with strong (naive) <u>independence</u> assumptions. A more descriptive term for the underlying probability model would be "<u>independent</u>feature model".

In simple terms, a naive Bayes classifier assumes that the presence (or absence) of a particular feature of a class is unrelated to the presence (or absence) of any other feature. For example, a fruit may be considered to be an apple if it is red, round, and about 4" in diameter. Even if these features depend on each other or upon the existence of the other features, a naive Bayes classifier considers all of these properties to independently contribute to the probability that this fruit is an apple.

Depending on the precise nature of the probability model, naive Bayes classifiers can be trained very efficiently in a <u>supervised learning</u> setting. In many practical applications, parameter estimation for naive Bayes models uses the method of <u>maximum likelihood</u>; in other words, one can work with the naive Bayes model without believing in <u>Bayesian probability</u> or using any Bayesian methods.

In spite of their naive design and apparently over-simplified assumptions, naive Bayes classifiers have worked quite well in many complex real-world situations. In 2004, analysis of the Bayesian classification problem has shown that there are some theoretical reasons for the apparently unreasonable <u>efficacy</u> of naive Bayes classifiers.<sup>[1]</sup> Still, a comprehensive comparison with other classification methods in 2006 showed that Bayes classification is outperformed by more current approaches, such as <u>boosted trees</u> or <u>random forests</u>.<sup>[2]</sup>

An advantage of the naive Bayes classifier is that it requires a small amount of training data to estimate the parameters (means and variances of the variables) necessary for classification. Because independent variables are assumed, only the variances of the variables for each class need to be determined and not the entirecovariance matix.

# **RESULTS:**



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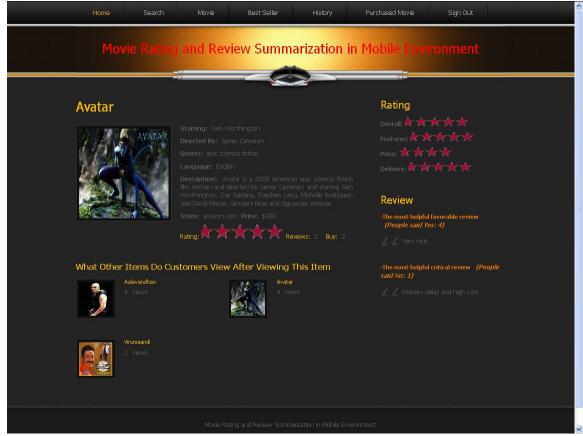
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		Festival of India.

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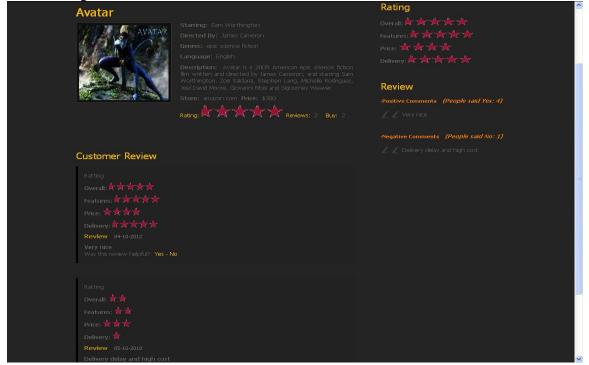
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	Directed By: James Cameron	
The second second	Genres: epic science fiction	
	Language: English	
	Description: Avatar is a 2009 American epic science fiction	
	film written and directed by James Cameron, and starring Sam Worthington, Zoe Saldana, Stephen Lang, Michelle Rodriquez,	
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# **Conclusions:**

We endeavored to modify a statistical stemming technique and came up with one whose results are comparable with a well known linguistic stemmer. Our N-gram stemmer is capable to deal any language for which N-gram corpus frequencies are available. As a language neutral approach is always preferable over a technique in which a linguistic knowledge or analysis is prerequisite, our results seem to be very exciting and promising.

It is really very important ant to identify the factors that improve the performances as well .factors that affect performance are the choice of training data, attribute or feature selection, representation of instances ,and the choice of the algorithm used. A decision to make equal number of each class that is positive and negative was being made to avoid biasing of classifier in favor of any particular class.

Attribute selection effected the accuracy that we can get from the classifier .too few attributes can cause the under fitting and too many attributes causing over fitting.

Under fitting s a situation in which the classifier is not biased enough to make prediction on some unseen instances. Over fitting is a situation where the train model is highly fitted to the training data that it basically fails while predicting new instances successfully. Thus it s good to find a balance between too few and too many attributes .unfortunately, there seems to be no as such simple way of finding the balanced number of attributes except by trail and error.

The other factors that affected the performance is the representation of instances .this basically includes whether to use presence /absence or count/frequency.it also includes whether to convert count into tf-idf so that terms importance is taken into account. That best result for multinomial naïve bayes can be obtained using this presence. Finally, the choice of machine learning algorithm for the task also affected the performance. Multinomial naïve bayes was found to outperform others in the task of sentiment analysis /classification task.

# **Future Work**

It is not possible to develop a system that meets all the requirements of the user. User requirements keep changing as the system is being used. Some of the future enhancements that can be done to this system are:

• As the technology emerges, it is possible to upgrade the system that can be adaptable to desired environment.

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