Dynamic Resource Allocation

Project report submitted in partial fulfillment of the requirement for the degree of Bachelor of Technology

in

Computer Science and Engineering/Information Technology

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Under the supervision of Dr. Monika Bharti (Assistant Professor, CSE)



Department of Computer Science & Engineering and Information Technology

Candidate's Declaration

I hereby declare that the work presented in this report entitled "Dynamic Resource Allocation " in partial fulfillment of the requirements for the award of the degree of Bachelor of Technology in Computer Science and Engineering/Information Technology submitted in the department of Computer Science & Engineering and Information Technology, Jaypee University of Information Technology Waknaghat is an authentic record of my own work carried out over a period from August 2023 to May 2023 under the supervision of (Dr. Monika Bharti) (Assistant Professor, CSE).I also authenticate that I have carried out the above mentioned project work under the proficiency stream, Cloud Computing. The matter embodied in the report has not been submitted for the award of any other degree or diploma.

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This is to certify that the above statement made by the candidate is true to the best of my knowledge.

(Supervisor Signature) Supervisor Name:Dr Monika Bharti Designation:Assistant Professor Department name:CSE Dated:26/04/2023

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LIST OF ACRONYMS AND ABBREVIATIONS

CC-Cloud Computing

Cloudsim:CloudSim is a framework for modeling and simulation

Task Scheduling Algorithms

VM-Virtual machine

Algo-Algortithm

Web App-Website based Application

LIST OF FIGURES

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ABSTRACT

Cloud computing manage a variety of virtualized resources which makes scheduling a critical component. In the cloud, a client may utilize several thousand virtualized assets for every task. Consequently, a manual scheduling is not a feasible solution. The basic idea behind task scheduling is to slate tasks for the minimization of time loss and maximize performance. Several research efforts have examined task scheduling in the past. This project presents a comprehensive survey of task scheduling strategies and the associated metrics suitable for cloud computing environments. Distinctive scheduling procedures are studied to discover which characteristics are to be included in a given system and which ones to be disregard. The project presents a review study of various task scheduling algorithms in cloud environment including RR,MaxMin,MinMin,FCFS,Selective,SJF. The algorithmhas been tested using CloudSim toolkit. It is recommended to use the algorithms with best make span time ,average waiting time,completion cost in order to avoid starvation and making it suitable for load balancing.

Keywords : CloudSim, FCFS, Genetic Algorithm, Round Robin Scheduling Algorithm, Virtual Machine Scheduling,SJF Algorithm,Maxmin Algorithm,Minmin Algorithm,Selective Algorithm.

Chapter 1 Introduction

1.1 Introduction

The term "Cloud" refers to network of linked computers that contains multiple unified computing resources. The development of cloud computing in recent years has made it possible to quickly assemble a network of geographically scattered data centres for the provision of high-quality and reliable services . Cloud computing is now a successful paradigm for providing computational capabilities on a "pay-per-utilize" basis. The IT industry is changing and becoming more uniform thanks to cloud computing. Cloud computing has many opportunities for growth while also posing numerous challenges to the evolution of traditional IT due to its expanding application and promotion. Cloud computing has recently become a popular IT(internet) strategy empowering users. It provide access to common reservoir Configurable resources that are available immediately and can be released with minimum management or help from cloud providers. This innovation has a number of benefits, including the ability to enhance time, cost, stack adjusting, and storage benefits in the marketplace. With the help of this invention, all applications can continue to operate on a virtual platform, and all resources are divided the VM. Each and application is unique independent.

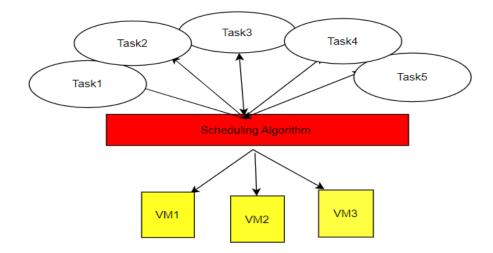


Fig 1.1 Scheduling Algorithms

Some parallel programmes display a decrease in CPU resource utilisation when parallelism is increased. If the jobs are not scheduled appropriately, performance degrades due to the cloud operations' enormous volume of data. Thus, in the sphere of cloud computing, the scheduling method is crucial. To plan and carry out the tasks with the highest assessed gain or advantages, a scheduling algorithm is used. Users use tens of thousands of resources, making it hard to assign each work by hand. The VM layers handle the scheduling tasks. In cloud computing, scheduling is crucial for allocating resources to each activity effectively and efficiently.

1.1.1 VM Scheduling

Processes or activities are planned according to the provided requirements and employed methods in a balanced scenario called scheduling. Cloud computing,Virtual Machine scheduling algorithms used for scheduling , Virtual servers requests to the actual machine of specific data centers in accordance with the needs satisfied by the requested resources. Schdeuling Algorithms works on the following rules.

- For VM find PM.
- Determine the allocation for VM
- Schedule the jobs for VM.

As was mentioned in the introduction section, a significant component of managing cloud resources cost-, energy-, and security-wise is cloud deployment. Three phases make up cloud provisioning: provisioning virtual machines, provisioning resources, and provisioning applications [7]. Here, provisioning for virtual machines must be the main concern.

Cloud deployment is also a key factor for managing resources in a costeffective, energy efficient, cloud provision consists of 3 step: (i) CM provision, (ii) resource (iii) APP provision. Here we need to focus on VM Allocation.

Task SchedulingAlgorithmsCategories:

- Instant Scheduling-A new task arrives, these scheduled directly on the VMs.
- BatchSchedulingTasks is combined batches before submission. These types are also called a mapping event.
- Static scheduling- regarded as being considerably simpler than dynamic scheduling. This is based on prior knowledge about the system's general condition.
- Dynamic Scheduling- Jobs are distributed according to the capacity of all the available VMs, taking into consideration the VM(s) current states and without knowing the system's overall condition beforehand.
- Preemptive scheduling-Each of the task can be interrupted while execution and resource to complete execution.
- NonPreemptiveScheduling-VM are not assign new tasks are scheduled tasks have finished running.

1.1.3 Task Scheduling Method

- A task is sent to a scheduler by a CC user. For resource information, a scheduler communicates with the cloud information system.
- The scheduler received the resource information from CIS.
- The scheduling algorithm performs its function by assigning tasks to the appropriate resources and submitting them to the successful resources.
- Through a cloud interface, the user obtains the resource's identifier and makes use of it.
- According to the timetable, the user delivers the resources.
- To control schedule scheduler receives periodically update the information the cloud condition.

1.1.4 Task Scheduling system in Cloud Computing

- There are 3 tiers in the cloud computing work scheduling system.
- A group of tasks send the cloud users that are necessary for exec. make up the first task level.
- The second scheduling level is in charge of assigning tasks to the appropriate resources to achieve the best resource utilization with the shortest lead time. The makespan time is the total time required complete the tasks from beginning end.
- The tasks shown in the accompanying figure are carried out at the third level by a group of people.

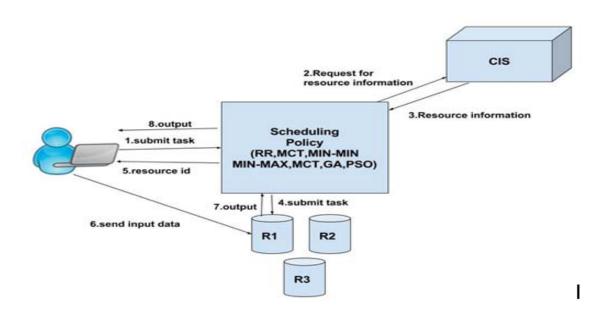


Fig 1.2 Task scheduling System

1.1.5 Types of Task Scheduling

I. Dynamic Task scheduling

The preceding information about framework applications and resources are not required for the dynamic load balancing. Shift the process from highload machines to low-load machines. Supports process slicing . There are various wellknown algos used in CC environments.

II. Static Load balancing

It needs previously specified details framework app,resources. A load balancing algorithm is referred to as static load balancing. distributes the backlog according to a set of fixed rules that are identified by characteristics of the currently active tasks.

III. Task scheduling algorithms advantages

- Control the QoS and performance of cloud computing.
- Control the CPU and the memory. .
- Improving the fairness of all the tasks.
- Scheduling tasks in real time basis.
- Improving load balancing.
- Achieving a high system throughput.
- Completed the tasks on time .

1.1.6 Scheduling Algorithms

I RR algortihm

The RR algo mainly idea of time slice . Time is divided several quantums. Each node receives a specific time Quantum, during the time, node performs all tasks . The RR is based sampling. Select loads in random manner where some servers are lightly loaded or some are heavy loads.

Pros:

- The suggested VM can be retrieved without any preprocessing steps.
- Equally distribute the load among the VMs.
- emphasizes equity among the assigned duties.
- Jobs are carried out sequentially, without ever waiting for the completion of a preceding job (starvation free).
- The scheduler won't wait for a VM to run out of processing power before moving on to the next one.
- It is founded on a basic principle.

Cons:

- Long tasks require extra time to execute.
- overloading
- Policies slicing are dependent on the size of the time slice, and in the case slice, many switching occur.

II. SJF Algorithm :

The SJF policy selects jobs that require shortest process time first. Shorter jobs are then run first. The main problem with SJF is the need to calculate the processing time for each job. If two jobs have the same processing time, SJF uses fcfs policy. This algo. also known as the shortest cpu burst time algo.

Pros :

• Decrease the average wait time since it decreases the wait time for quick jobs.

• increase the avg. waiting time of heavy tasks.

Cons:

• Starvation for heavy tasks

III. FCFS:

FCFS algorithm is simplest scheduling algorithm. When job come the less executiontime job is done fist.FIFO queues are used to implement FCFS policies. Pros:

- Simple
- Easy to understand

Cons:

- Nonpreemptive.
- Less execution job is done so starvation.

IV. Min-Min Algorithm:

The Min-min method ensures that each task is completed in the shortest amount of time by allocating each task to a virtual computer known as a resource. Determine the duration of each job's execution on each resource that is available. There are two steps in the Minmin algorithm. The work with the quickest execution time is determined in first stage. Later on, the job with the fastest execution time is chosen. Then, algorithms assign resources to projects. Up till all jobs are mapped, this process is repeated.

Pros:

- Less makespan.
- No pre-checking for machines Increase the throughput.

Cons:

• The heavy tasks have to wait, Unbalanced load

V. Maxmin Algorithm

The maxmin method maps each job to the resource VM in such a way that each job can be finished in reasonable amount of time. Every job's execution time on each resource that is available is calculated. The assignment that requires the most time to complete is then chosen. The maxmin algorithms require less waiting time for large jobs.

Pros:

- Avg waiting less.
- Util. Increases.
- Min response time.

Cons:

- selects big tasks.
- lead to starvation.

VII. Selective Algorthim

We developed an algorithm that distributes tasks to the best computers in a way that gives sufficient performance to both cloud customers and providers in order to meet our objectives of lowering the total makespan of jobs on machines and providing improved quality of service. Based on specific criteria, the method is intended to select between the minmin scheduling algorithm and the maxmin algorithm implemented.

Pros:

- Less waiting time.
- local optimization
- Response time is minimized.

Cons:

- Not effective in load balancing
- Small makespan time.

1.2 Problem Statement

In order to minimise the execution time, or is MakeSpan time and resource cost the..total cost, the major challenge of task scheduling is to determine the optimal mapping of n independent tasks into m heterogenous resources. Maximising the resource utilization. Main goal of task schedul. creating map f to allocate (T), (R). i.e f:T->R.

- Tasks are independent
- The available resources (VMs) are heterogenous.

The project goal to min. the task execution time and execution cost as well increase resource utilization.

Main challenges are :

Resource Distribution Internet-based, cost-effective deployment of cloud resources to multiple applications.

The process of providing a service provider's resources to cloud users while assuring service quality is known as resource provisioning. This procedure is based on the service level agreement (SLA). It is classified as either Dynamic Resource Provisioning or Static Resource Provisioning.

Resource mapping: It is the alignment of resources needed by cloud customers with resources offered by a service provider.

Resource discovery and selection: The process of discovering all resources present in the system, gathering the current state of the resources, and deciding which target resources to select based on the information gathered from the discovery.

Resource adaption: is the capability of this system to dynamically adjust reources to meet the user requirements.

I. Enumeration Method

If all potential solutions are listed and weighed against one another, the best one for an optimization problem can be chosen. In the worst situation, exact enumerative algorithms have exponential time complexity. However, for a few weakly defined NP-hard problems, when the

When number of instance is small enough, it be solved using a pseudopolynomial method, the time complexity of which is constrained the input size and the problem's maximum number and their polynomial form. In addition, there is a different type of enumeration known as implicit enumeration, which examines every potential answer without explicitly naming them all. A useful implicit enumeration technique to handle combinational optimization issues is dynamic programming. It breaks down a problem into several stages, and The decisions that must be made at each stage have an effect on those that must be made at later stages. Since the number of stored decisions grows exponentially with the number of subproblems, exponential complexity is the worst type of complexity for dynamic programming algorithms.

II. Heuristic Method

Due to the fact that onlyfew particular situations of NP-hard problems have absolutely solution algorithms polynomial time, exhaustive enumeration is not practical for scheduling issues. We frequently settle for less-than-ideal solutions in order to balance accuracy and time since it's good practice. A heuristic is a suboptimal algorithm that locates moderately effective answers quickly. It does not ensure the optimal solution, but improves a candidate solution in relation to a specific criterion quality. In order to assess the precision of heuristic algorithms, approximation rate rH(e) is introduced .

III. Relaxation Method

Relaxing some of the original problem's constraints is another viable approach to solving NP-hard issues. The answer to the new relaxed problem might simple to find and have close approximation the solution to the problem. The typical easing includes:

• Assume that instead of real numbers, all the constituents in one instance are natural numbers.

• Assume that the value of one particular element doesn't change at all.

• Assume that two linked elements have equal values rather than that one of them is constrained by the other.

- Assume that an element's value is unit rather than being chosen at random.
- Assume that an element's type is predetermined rather than random.

1.3 Objective

1.3.1 Resource provisioning

When the people request for resources is approved by cloud service provider, a resource allocation technique is used to create and assign that user a practical number of virtual machines (VMs) based on demand. In addition, resource provisioning is in charge of matching incoming workloads or applications (cloudlets) to resources, user requirements based on QoS standards, and SLA agreements (VMs). In other words, it is important to provide incoming tasks and applications with the fewest resources possible while still ensuring a high level of quality of service.

1.3.2 Scheduling resources

The scheduling involves examining tnecessary QoS characteristics in order to decide which activity should be carried out.

- Selecting the most appropriate VM.
- ensuring the fulfillment of QoS.

The RPS handle the task scheduling , and aims to provision the VMs to users.

• understanding the constomer demands.

1.4 Methodology

There aren't many sources of canonical inputs or benchmarks to conduct comparison assessments because cloud computing is a relatively new sector. Running a benchmark additionally necessitates testing in a repeatable, dependable, and scalable environment, which is not feasible in the real world given the variety of cloud service providers and various planning policies. is attainable. and surroundings.

So, to get an overall software framework for modeling a cloud computing environment, and to run the tests, We employ a cloud simulator by the name of CloudSim. The flexible and expandable simulation framework that CloudSim offers makes it possible to model, simulate, and experiment with novel cloud computing infrastructure and application services with ease. With the aid of CloudSim, developers can concentrate on the specific system design issue wish to investigate without having to worry about the nitty-gritty specifics of cloud-based infrastructure and services.

We propose to use CloudSim to compare the Execution time . cost, makespan time to the other scheduling algorithms common inputs.

Algorithm Pseduo Code:-

To this end, use two registers to store the total burst time of all requests in the RQ, and use Areg to divide the value in SReg by the no. of task in the RQ. Stores the average burst time by column.

Inputs -SReg, AReg, Tn (task n), BT, TQ, Ready Queue

1.Begin

- 2. New requests Task arrives, tasks enter RQ
- 3. Update Sreg and Areg requests.
- 4. Task are loaded from RQ into VM queue be executed.
- 5. While (Ready Queue! = NULL) do
- 6. Ready queues T
- 7. Update Sreg&Areg
- 8. load T
- 9. if(readyQueue!=NULL) do
- 10. tq=BT(T)
- 11 incr. SReg& AReg

else

- 12. tq=avg(BT of all requests).
- 13. Sreg
- 14. // VM executes T by TQ Time
- 15. if(T terminated) then
- 16. Update Sreg & Areg
- else Retrun T
- 18. Burst Time(BT)
- 19. Update AReg
- 20. end if

1.5 Organization

The content present in the prject report is designed into. five chapters firstly the introduction chapter , chpater 2 describes literature survey, where different reseearch papers related to this project work have been included on task scheduling, resource allocation. Mostly explaned resource allocation , vm scheduling , task scheduling algorithms and considering the factors affecting task scheduling.

Chapter 3 summarizes the system design where various analytical and devlopment analysis is explained along with the design and algorithm of the mdel devepment. Model development technique is explained along with the design and agorithms of the model development. Model development technique has been explained using cloudsim simualtions.Web App framework development cycle.

Chapter 4 provides an account of the performance analysis where we have mentioned the most suitable algorithms for this prject after cmoparing the agrithms the basis of Makespantime, total execution cost, avg waiting time.

Chapter 5 reprents a brief summary of conclusion, and future work based on the research done during the implementation of the project.

In the end References is added where all the research papers are mentioned that were needed fot the better implementation of the algorithms.

Chapter 2-Literature survey

The best scheduling policy in the hydrid cloud model was proposed by R. Vanmechelen, and J. Broeckhove in 2010, according to a review of the literature. Both private and public clouds are included in the hybrid cloud. Private and public clouds make up hybrid. The situation of the issue is that, during periods of peak load, some workloads must be transferred from private cloud to public cloud since there aren't enough resources in the former to handle the tasks given by users. Both a deadline and QoS specifications place restrictions on these workloads. In this situation, a decision-making process is required to choose which workloads to which cloud provider, order to maximize amount of resources used in the internal data center and reduce operating costs.

In 2011, S. Sindhu and Saswati Mukherjee suggested algorithms for task scheduling in cloud computingon task processing requirements and resource computing power. The first algorithm, called (L CFPE), allocates longer cloudlets (tasks) than processing elements (PEs) with higher values and reduces the Makespan (time it takes to complete all tasks. total time). A second algorithm called Shortest Cloudlet Fastestb Processing

Dandhwani,Vanita,andVipul presented a K-means- task scheduling algorithm. The idea of the proposed algorithm is to use k-mean clustering technique to create clusters of tasks and assign clusters to VMs according to VM capacity. Results showed that the execution time and makespan of the were reduced, improving the overall system.

A allocation (PACO) method was proposed for work scheduling in the cloud system by Ruonan Lin and Li Qiang in 2016. The template size and

the enhanced Ant Colony Optimization (ACO) method are used into this technique to individually plan activities. The simulated software tests the proposed algorithm and finds it to be rather efficient. The experiments show that the productivity of PACO Scheduling will improve the task.

2015, Moradbeiky A. Bardsiri conducted research using a task scheduling algorithm based on Cuckoo Optimization. In this algorithm, the bird's nest simulates processing unit (a virtual machine) and the egg is the task.

The parallel GA based approach for prioritizing the tasks when scheduling them in cloud systems was presented by Mehran Ashouraei and colleagues in 2018 (Mehran Ashouraei, 2018). This approach aims to use resources efficiently and reduce resource waste. In order to eliminate work failures, this method involves increasing load balance choosing qualified resources for urgent task.

Element (SCFP) does the opposite. SCFP maps short cloudlet process. Using this algorithm minimizes lead times and avoids running out of longer orders. They suggested that future work should use heuristics to test more algorithms and consider task priorities. It is used to optimize scheduling under the cloud . ACO resembles the behavior of ant colonies, in which ants move in random directions in search of food sources. In task scheduling, tasks such as ants and virtual machines mimic food sources.

CHAPTER 3: SYSTEM DEVELOPMENT

3.1 Software Specifications

All algorithms were written in Java 8 as .java files.Java was chosen because it is a powerful compiler. Java supports both Oops execution and debugging at the same time.Furthermore, Java is platform independent, unlike other languages like C++, C, which compile onto platform but Java is write-once, run-anywhere language. Java differs from other platform it is software based on the other platforms based software, ithas two components.

Java programming is platform-independent. Linux, Mac OS, Sun Solaris, Windows, etc. A compiler compiles Java code and turns it into bytecode. This bytecode can run on various platforms, making it platform independent code. Run anywhere and write once.

The general purpose, interactive, object-oriented, and high-level programming language python is particularly liked by global users, Python is dynamically and garbage collected.

Docker system software letting users to deploy, build, testing the web applications, Software is packaged by docker into standardised units called containers that contain all of the necessary libraries, system tools, code, and runtime for the software to function.

3.1.1 Cloudsim Toolkit

Cloudism is open source JDK tool originally distributed ,laboratory in the melbourne university. Cloudsim was used to simul the components of a cloud computing system, evaluate resource provisioning policies, test various Cloud Computing deployments, and generate a mix of workload requirements distribution. , to measure the performance of CC configure. and test scenario. These capabilities provided by CloudSim, Cloud sim developers to efficiently address and operate several key issues related to cloud computing and ultimately

develops best results.

The complex real-world CC environment is modelled using the Cloudsim framework, which also simulates its behavior. With the help of this tool, a model is made up of many elements, including datacenters, hosts, service brokers, scheduling, and allocation policies. Cloud computing will develop.

To start using your CC system, you must have at least one data center listed in the Cloud Information Service Registry (CIS). Each data center creates a number of hosts. After receiving the list of cloudlets and virtual machines that have been submitted, the Data Center Broker applies allocation policy assign servers to VMs.

The CloudSim software is impl. using a layer design as shown in below Figure3.3 The first layer is simjava, which provides core functionalityto the top cloud layers, such as creating cloud systems and communicating between them. increase, Components simulation clock management, and queuing and event handling.

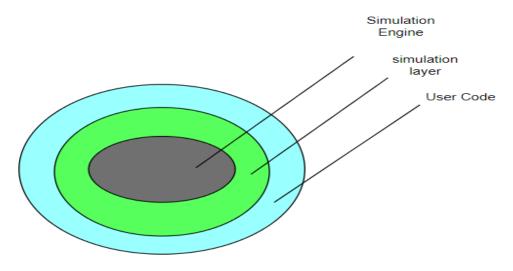


Fig 3.3 CloudSim layers

Eclipse

An integrated development environment called Eclipse is employed in computer programming. The environment can be customized using an extendable plug-in system and a base workspace. It is the second-most widely used IDE for developing Java.

3.2 Hardware Specification

Algorithms that are platform dependent have been implemented. They may be run on any of the Windows,Mac systems. We wrote the coding for our initial prototype in Eclipse. The GPU was utilized to build the model.

3.3 Input

To ensure the relevance of a simulation-based evaluation, investigations must be conducted utilizing workload are from real world. Input data for CPU utl. for over 100 VMs from servers .20 tasks, 20 cludlets The utilization measurement interval is 5 minutes, each tracked file has, so represents approximately 24,44 hours of VM CPU utilization.

For Task scheduling simulator the input are given by user like:

- execution time
- arrival time
- period
- deadline

3.4 Output

It is noted that SLA violations and energy usage occur for activities that range in size from sets of 10 to 100. To determine the true statistical performance difference between the methods, simulation is performed and result are drawn that compare the results from the suggested method and existing scheduling algorithms. the make span time is considered in the performance and comparison of schedling algorithms.

The simulator gives output result in the form of graph.

3.5 Algorithm Designs

Tasks T = to, t1...t2...tN-1, and no.of tasks is N=|T| whereas the no of hosts is m=|H| for the set ... of hosts H = to, h0, h1...hn-1, The avg. load of the VM that run a host system is referred to ... as the host machine's load.

Step 1: Set the host resource $H=\{h0,h1...hm-1\}$ and sort in asc. order of the processing power.

Step 2: Considering the VM, choose the host resource (that can handle the least amount of load while providing the necessary resources. If the host is already present, create the virtual machine and allocate the appropriate resources for it, then update the resources that are now accessible. Move Task t1 to the end of the task queue in the absence of a host and wait for the next scheduling.

Step 3: If the Task t i's resource requirements rise, determine whether the host that Task t i's virtual machine runs on can accommodate the extra resources needed, if so, allocate those resources to the host, modify the virtual machine's configuration before updating the host's resource availability.

If not, the virtual machine is transferred to the host that experiences the least load and has the additional resources required to run continuously.

Step 4: If resourcess requires task t1 reduce release the overlaod resources so VM occupied and update the host.

Step 5: If taks t1 has been completed then destroys then destroy the virtual machine of task t1 realse the occupied resource for the other unfinished tasks.

Step 6: Find the SDV. of each host's load. Select one virtual serer from a host and relocate it to host with the lowest load if one host has a significantly higher load than the others. If a host is under very low stress, move all virtual machines (VMs) from that host to another host.

Step 7:

Determine lad on every host and its std.deviation. If one host has a much larger load than the others, choose a virtual machine from that host and move it to the host with the lowest load. Move all virtual machines (VMs) from a given host if it is experiencing very little load to another host.

3.6 Creating Web Application

With the increase of technical aspect of life, we aim to create an scheduling simulator for task scheduling,

with one platform for scheduling simulation for all scheduling algorithms, for giving an idea of how simulation works.

So to make our application user ready we have created an interface using technologies like HTML,

CSS,Django.

3.6.1 HTML

HTML stands for Hyper Text Markup Language. It will act as the skeleton for Web page and its design.

3.6.2 CSS

It stands for Cascading Style Sheet. It is mainly used to beautify our web page.

3.6.3 Django

A Python-based web framework called Django enables you to easily build effective online apps. Django has built-in functionality for everything, including the Django Admin Interface and the default SQLlite3 database, it is often known as the batteries included framework. It is highly scalable, easy to develop web app and learn.

Flowchart

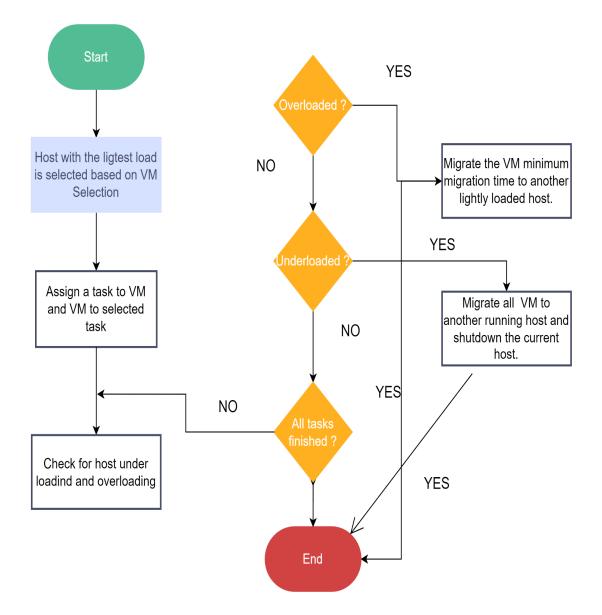


Fig 3.4 System Design Flowchart

Chapter 4: Performance and Analysis

4.1 Algortihms Results and Analysis

1.FCFS Algortihms:

atacenter_1 atacenter_1 roker_0 is : imulation c imulation c ==================================0	9 is shutti shutting do ompleted. ompleted.	ng down wn									
loudlet ID	STATUS	Data center ID	VM ID	Time	Start Time	Finish Time	Wa	itingTime	Completi	onTime	Cost
38	SUCCESS	03	03	00	00.1	00.1	00	00	00		0000
06	SUCCESS	07	07	00	00.1	00.1	00	00	00		
10	SUCCESS	08	08	00	00.1	00.1	00	00	00		
02	SUCCESS	09	09	00	00.1	00.1	00	00	00		
07	SUCCESS	10	10	0.0	00.1	00.1	00	00	00		
15	SUCCESS	11	11	00	00.1	00.1	00	00	00		
11	SUCCESS	12	12	00	00.1	00.1	00	00	00		
08	SUCCESS	13	13	00	00.1	00.1	00	00	00		
04	SUCCESS	14	14	00	00.1	00.1	00	00	00		
05	SUCCESS	15	15	00	00.1	00.1	00	00	00		
12	SUCCESS	16	16	00	00.1	00.1	00	00	00		
42	SUCCESS	17	17	00	00.1	00.1	00	00	00		
28	SUCCESS	18	18	00	00.1	00.1	00	00	00		
00	SUCCESS	19	19	00	00.1	00.1	00	00	00		
14	SUCCESS	20	20	00	00.1	00.1	00	00	00		
41	SUCCESS	21	21	00	00.1	00.1	00	00	00		
39	SUCCESS	03	03	00.11	00.1	00.21		00	00.11	00.33	
36	SUCCESS	07	07	00.11	00.1	00.21		00	00.11	00.33	
20	SUCCESS	08	08	00.11	00.1	00.21		00	00.11	00.33	
32	SUCCESS	09	09	00.11	00.1	00.21		00	00.11	00.33	
17	SUCCESS	10	10	00.11	00.1	00.21		00	00.11	00.33	
23 18	SUCCESS	11	11 12	00.11	00.1	00.21		00	00.11	00.33	
25	SUCCESS SUCCESS	12 13	12	00.11 00.11	00.1 00.1	00.21 00.21		00	00.11 00.11	00.33 00.33	
25 44	SUCCESS	13	13	00.11	00.1	00.21		00	00.11	00.33	
44 30	SUCCESS	15	15	00.11	00.1	00.21		00	00.11	00.33	
01	SUCCESS	19	10	00.11	00.1	00.21		00	00.11	00.33	
16	SUCCESS	20	20	00.11	00.1	00.21		00	00.11	00.33	
48	SUCCESS	09	09	00.11	00.21	00.32		00.11	00.11	00.33	22
22	SUCCESS	10	10	00.11	00.21	00.32		00.11	00.22	00.	
47	SUCCESS	10	10	00.11	00.21	00.32		00.11	00.22	00.	
26	SUCCESS	12	12	00.11	00.21	00.32		00.11	00.22	00.	

39	SUCCESS	03	03	00.11	00.1	00.21	00	00.11	00.33	
36	SUCCESS	07	07	00.11	00.1	00.21	00	00.11	00.33	
20	SUCCESS	08	08	00.11	00.1	00.21	00	00.11	00.33	
32	SUCCESS	09	09	00.11	00.1	00.21	00	00.11	00.33	
17	SUCCESS	10	10	00.11	00.1	00.21	00	00.11	00.33	
23	SUCCESS	11	11	00.11	00.1	00.21	00	00.11	00.33	
18	SUCCESS	12	12	00.11	00.1	00.21	00	00.11	00.33	
25	SUCCESS	13	13	00.11	00.1	00.21	00	00.11	00.33	
44	SUCCESS	15	15	00.11	00.1	00.21	00	00.11	00.33	
30	SUCCESS	16	16	00.11	00.1	00.21	00	00.11	00.33	
01	SUCCESS	10	10	00.11	00.1	00.21	00	00.11	00.33	
16	SUCCESS	20	20	00.11	00.1	00.21	00	00.11	00.33	
48	SUCCESS	09	09	00.11	00.21	00.32	00.11	00.11	00.33	
48 22	SUCCESS	10	10	00.11	00.21	00.32	00.11	00.22	00.33	
47	SUCCESS	10	10	00.11	00.21	00.32	00.11	00.22	00.33	
26	SUCCESS	11	12	00.11	00.21	00.32	00.11	00.22	00.33	
20		12	12	00.11	00.21	00.32	00.11	00.22	00.33	
31	SUCCESS SUCCESS	13 16	13	00.11	00.21	00.32	00.11	00.22	00.33	
03			10	00.11		00.32	00.11	00.22	00.33	
	SUCCESS	19			00.21					
24	SUCCESS	20	20	00.11	00.21	00.32	00.11	00.22	00.33	
29	SUCCESS	12	12	00.11	00.32	00.43	00.22	00.33	00.33	
35	SUCCESS	13	13	00.11	00.32	00.43	00.22	00.33	00.33	
34	SUCCESS	16	16	00.11	00.32	00.43	00.22	00.33	00.33	
46	SUCCESS	19	19	00.11	00.32	00.43	00.22	00.33	00.33	
40	SUCCESS	20	20	00.11	00.32	00.43	00.22	00.33	00.33	
33	SUCCESS	12	12	00.11	00.43	00.54	00.33	00.44	00.33	
43	SUCCESS	13	13	00.11	00.43	00.54	00.33	00.44	00.33	
49	SUCCESS	16	16	00.11	00.43	00.54	00.33	00.44	00.33	
45	SUCCESS	13	13	00.11	00.54	00.65	00.44	00.55	00.33	
09	SUCCESS	04	04	522.55	00.1	522.65	00	522.55	1567.64	ł
21	SUCCESS	02	02	637.97	00.1	638.07	00	637.97	1913.9	
37	SUCCESS	02	02	00.11	638.07	638.18	637.97).33
13	SUCCESS	04	04	715.46	522.65	1238.11	522.		238.01	2146.3
19	SUCCESS	04	04	666.37	1238.11	1904.47	123	3.01	1904.38	1999
	sing FCFS: 1221									
	letion Time: 49									
	: 7636.9409999		st: 152.738819	99999998						
	g Time: 48.0388									
g.cloudbu	us.cloudsim.exa	mples.FCFS_S	cheduler finis	hed!						

Makespan using FCFS: 1221.2232289234394

Total Completion Time: 4947.59 Avg Completion Time: 98.9518

Total Cost : 7636.940999999999 Avg cost: 152.73881999999998

Avg Waiting Time: 48.0388600000000

2.MinMin algorithm

Simulation completed.

Cloudlet ID	STATUS	Data center ID		Time		Finish Time
48	SUCCESS	5	5	0	0.1	0.1
0	SUCCESS	7	7	0		0.1
24	SUCCESS	8	8	0	0.1	0.1
23	SUCCESS	9	9	0	0.1	0.1
10	SUCCESS	10	10	0	0.1	0.1
27	SUCCESS	11	11	0	0.1	0.1
8	SUCCESS	12	12	0	0.1	0.1
12	SUCCESS	13	13	0	0.1	0.1
15	SUCCESS	14	14	0	0.1	0.1
3	SUCCESS	15	15	0	0.1	0.1
29	SUCCESS	16	16	0	0.1	0.1
21	SUCCESS	17	17	0	0.1	0.1
30	SUCCESS	18	18	0	0.1	0.1
1	SUCCESS	19	19	0	0.1	0.1
20	SUCCESS	20	20	0	0.1	0.1
2	SUCCESS	21	21	0	0.1	0.1
19	SUCCESS	7	7	0.11	0.1	0.21
13	SUCCESS	10	10	0.11	0.1	0.21
36	SUCCESS	11	11	0.11	0.1	0.21
25	SUCCESS	12	12	0.11	0.1	0.21
28	SUCCESS	13	13	0.11	0.1	0.21
41	SUCCESS	14	14	0.11	0.1	0.21
31	SUCCESS	16	16	0.11	0.1	0.21
43	SUCCESS	17	17	0.11	0.1	0.21
11	SUCCESS	19	19	0.11	0.1	0.21
7	SUCCESS	21	21	0.11	0.1	0.21
32	SUCCESS	7	7	0.11	0.21	0.32
17	SUCCESS	10	10	0.11	0.21	0.32
46	SUCCESS	11	11	0.11	0.21	0.32
39	SUCCESS	12	12	0.11	0.21	0.32
35	SUCCESS	13	13	0.11	0.21	0.32
14	SUCCESS	19	19	0.11	0.21	0.32
34	SUCCESS	10	10	0.11	0.32	0.43
45	SUCCESS	12	12	0.11	0.32	0.43
22	SUCCESS	19	19	0.11	0.32	0.43
4						

13	SUCCESS	10	10	0.11	0.1	0.21
36	SUCCESS	11	11	0.11	0.1	0.21
25	SUCCESS	12	12	0.11	0.1	0.21
28	SUCCESS	13	13	0.11	0.1	0.21
41	SUCCESS	14	14	0.11	0.1	0.21
31	SUCCESS	16	16	0.11	0.1	0.21
43	SUCCESS	17	17	0.11	0.1	0.21
11	SUCCESS	19	19	0.11	0.1	0.21
7	SUCCESS	21	21	0.11	0.1	0.21
32	SUCCESS	7	7	0.11	0.21	0.32
17	SUCCESS	10	10	0.11	0.21	0.32
46	SUCCESS	11	11	0.11	0.21	0.32
39	SUCCESS	12	12	0.11	0.21	0.32
35	SUCCESS	13	13	0.11	0.21	0.32
14	SUCCESS	19	19	0.11	0.21	0.32
34	SUCCESS	10	10	0.11	0.32	0.43
45	SUCCESS	12	12	0.11	0.32	0.43
22	SUCCESS	19	19	0.11	0.32	0.43
37	SUCCESS	10	10	0.11	0.43	0.54
47	SUCCESS	12	12	0.11	0.43	0.54
42	SUCCESS	19	19	0.11	0.43	0.54
49	SUCCESS	10	10	0.11	0.54	0.65
4	SUCCESS	4	4	122.08	0.1	122.18
33	SUCCESS	4	4	0.11	122.18	122.29
40	SUCCESS	4	4	0.11	122.29	122.4
44	SUCCESS	4	4	0.11	122.4	122.51
9	SUCCESS	6	6	316.06	0.1	316.16
6	SUCCESS	3	3	605.83	0.1	605.93
5	SUCCESS	2	2	779.19	0.1	779.29
18	SUCCESS	3	3	278.9	605.93	884.83
16	SUCCESS	6	6	742.39	316.16	1058.56
26	SUCCESS	3	3	943.01	884.83	1827.84
38	SUCCESS	3	3	0.11	1827.84	1827.95
Makespan (using Minmin: 13	208.4114221	377815			
rotal Com	pletion Time: 7	794.141 Avg	Completion 7	Time: 155.882	81999999998	
rotal Cost	t : 11371.32900	0000002 Avg	cost: 227.42	265800000003		
	ng Time: 80.073					
org.cloud	bus.cloudsim.exa	amples.MinM	in finished!			
4						

Makespan using Minmin: 1208.4114221377815

Total Completion Time: 7794.141 Avg Completion Time: 155.88281999999998 Total Cost : 11371.32900000002 Avg cost: 227.4265800000003 Avg Waiting Time: 80.07396

3.MaxMin Algorithm

Simulation	completed.

(OUTPUT =====	=====				
Cloudlet ID	STATUS	Data center ID	VM ID	Time	Start Time	e Finish Time
45	SUCCESS	4	4	0	0.1	0.1
35	SUCCESS	5	5	0	0.1	0.1
12	SUCCESS	7	7	0	0.1	0.1
2	SUCCESS	8	8	0	0.1	0.1
4	SUCCESS	9	9	0	0.1	0.1
1	SUCCESS	10	10	0	0.1	0.1
17	SUCCESS	11	11	0	0.1	0.1
9	SUCCESS	12	12	0	0.1	0.1
19	SUCCESS	13	13	0	0.1	0.1
11	SUCCESS	14	14	0	0.1	0.1
7	SUCCESS	15	15	0	0.1	0.1
10	SUCCESS	16	16	0	0.1	0.1
8	SUCCESS	17	17	0	0.1	0.1
28	SUCCESS	18	18	0	0.1	0.1
14	SUCCESS	19	19	0	0.1	0.1
0	SUCCESS	20	20	0	0.1	0.1
15	SUCCESS	21	21	0	0.1	0.1
31	SUCCESS	7	7	0.11	0.1	0.21
26	SUCCESS	8	8	0.11	0.1	0.21
16	SUCCESS	10	10	0.11	0.1	0.21
22	SUCCESS	11	11	0.11	0.1	0.21
21	SUCCESS	12	12	0.11	0.1	0.21
29	SUCCESS	13	13	0.11	0.1	0.21
30	SUCCESS	14	14	0.11	0.1	0.21
24	SUCCESS	16	16	0.11	0.1	0.21
23	SUCCESS	19	19	0.11	0.1	0.21
6	SUCCESS	20	20	0.11	0.1	0.21
37	SUCCESS	10	10	0.11	0.21	0.32
33	SUCCESS	11	11	0.11	0.21	0.32
25	SUCCESS	12	12	0.11	0.21	0.32
34	SUCCESS	13	13	0.11	0.21	0.32
49	SUCCESS	16	16	0.11	0.21	0.32
39	SUCCESS	19	19	0.11	0.21	0.32
13	SUCCESS	20	20	0.11	0.21	0.32
38	SUCCESS	10	10	0.11	0.32	0.43
4						

.

- -

<terminated></terminated>	MaxMin [Java Applicat	ion] C:\Users\Shu	ubham\Desktop\clo	oud simsetup\eclip	ose-java-2022-06-R	-win32-x86_64\eclipse\			
31	SUCCESS	7	7	0.11	0.1	0.21			
26	SUCCESS	8	8	0.11	0.1	0.21			
16	SUCCESS	10	10	0.11	0.1	0.21			
22	SUCCESS	11	11	0.11	0.1	0.21			
21	SUCCESS	12	12	0.11	0.1	0.21			
29	SUCCESS	13	13	0.11	0.1	0.21			
30	SUCCESS	14	14	0.11	0.1	0.21			
24	SUCCESS	16	16	0.11	0.1	0.21			
23	SUCCESS	19	19	0.11	0.1	0.21			
6	SUCCESS	20	20	0.11	0.1	0.21			
37	SUCCESS	10	10	0.11	0.21	0.32			
33	SUCCESS	11	11	0.11	0.21	0.32			
25	SUCCESS	12	12	0.11	0.21	0.32			
34	SUCCESS	13	13	0.11	0.21	0.32			
49	SUCCESS	16	16	0.11	0.21	0.32			
39	SUCCESS	19	19	0.11	0.21	0.32			
13	SUCCESS	20	20	0.11	0.21	0.32			
38	SUCCESS	10	10	0.11	0.32	0.43			
32	SUCCESS	12	12	0.11	0.32	0.43			
44	SUCCESS	20	20	0.11	0.32	0.43			
42	SUCCESS	10	10	0.11	0.43	0.54			
36	SUCCESS	12	12	0.11	0.43	0.54			
46	SUCCESS	20	20	0.11	0.43	0.54			
48	SUCCESS	20	20	0.11	0.54	0.65			
27	SUCCESS	6	6	184.73	0.1	184.83			
41	SUCCESS	6	6	0.11	184.83	184.94			
18	SUCCESS	3	3	278.9	0.1	279			
40	SUCCESS	3	3	0.11	279	279.11			
43	SUCCESS	3	3	0.11	279.11	279.22			
3	SUCCESS	2	2	663.04	0.1	663.14			
5	SUCCESS	2	2	779.19	663.14	1442.33			
20	SUCCESS	2	2	388.4	1442.33	1830.73			
47	SUCCESS	2	2	0.11	1830.73	1830.84			
Total Com Total Cos Avg Waitin	Makespan using Maxmin: 1354.5280235668106 Total Completion Time: 6978.76200000001 Avg Completion Time: 139.57524 Total Cost : 6892.053000000001 Avg cost: 137.84106000000003 Avg Waiting Time: 93.62822 org.cloudbus.cloudsim.examples.Maxmin finished!								
org.croud	Sus.crouusin.ex	ampres.MdXIII.	in rinished:						
•									

Makespan using Maxmin: 1354.5280235668106

Total Completion Time: 6978.76200000001 Avg Completion Time: 139.57524

Total Cost : 6892.05300000001 Avg cost: 137.8410600000003

Avg Waiting Time: 93.62822

4. Round Robin Algorithm

<terminated>RoundRobinScheduler [Java Application] C\Users\Shubham\Desktop\cloud simsetup\eclipse-java-2022-06-R-win32-x86_64\eclipse\plugins\org.eclipse.justj.openjdk.hotspot.jre.full.win32.x86_64,17.0.3.v20220</terminated>
Simulation completed.

====== OU Cloudlet ID	STATUS	Data center ID	VM ID	Time	Start Time	Finish Time	Waiting	Time Comple	tionTime Co
.ioudiet iD 06	STATUS	07	07	00	00.1	00.1	00 Waiting	11me compie 00 00	CTOULTHUE CO
28	SUCCESS	08	08	00	00.1	00.1	00	00 00	
10	SUCCESS	09	09	00	00.1	00.1	00	00 00	
04	SUCCESS	10	10	00	00.1	00.1	00	00 00	
13	SUCCESS	10	10	00	00.1	00.1	00	00 00	
00	SUCCESS	12	12	00	00.1	00.1	00	00 00	
46	SUCCESS	13	13	00	00.1	00.1	00	00 00	
02	SUCCESS	14	14	00	00.1	00.1	00	00 00	
03	SUCCESS	15	15	00	00.1	00.1	00	00 00	
31	SUCCESS	16	16	00	00.1	00.1	00	00 00	
09	SUCCESS	17	17	00	00.1	00.1	00	00 00	
38	SUCCESS	18	18	00	00.1	00.1	00	00 00	
01	SUCCESS	19	19	00	00.1	00.1	00	00 00	
18	SUCCESS	20	20	00	00.1	00.1	00	00 00	
15	SUCCESS	21	21	00	00.1	00.1	00	00 00	
25	SUCCESS	07	07	00.11	00.1	00.21	00	00.11	00.33
21	SUCCESS	09	09	00.11	00.1	00.21	00	00.11	00.33
33	SUCCESS	11	11	00.11	00.1	00.21	00	00.11	00.33
24	SUCCESS	12	12	00.11	00.1	00.21	00	00.11	00.33
19	SUCCESS	14	14	00.11	00.1	00.21	00	00.11	00.33
08	SUCCESS	15	15	00.11	00.1	00.21	00	00.11	00.33
42	SUCCESS	16	16	00.11	00.1	00.21	00	00.11	00.33
11	SUCCESS	17	17	00.11	00.1	00.21	00	00.11	00.33
20	SUCCESS	19	19	00.11	00.1	00.21	00	00.11	00.33
27	SUCCESS	20	20	00.11	00.1	00.21	00	00.11	00.33
16	SUCCESS	21	21	00.11	00.1	00.21	00	00.11	00.33
40	SUCCESS	07	07	00.11	00.21	00.32	00.1	1 00.22	00.33
23	SUCCESS	09	09	00.11	00.21	00.32	00.1	1 00.22	00.33
47	SUCCESS	14	14	00.11	00.21	00.32	00.1		
30	SUCCESS	15	15	00.11	00.21	00.32	00.1		
45	SUCCESS	17	17	00.11	00.21	00.32	00.1	1 00.22	00.33
26	SUCCESS	19	19	00.11	00.21	00.32	00.1	1 00.22	00.33
39	SUCCESS	20	20	00.11	00.21	00.32	00.1		
22	SUCCESS	21	21	00.11	00.21	00.32	00.1		
35	SUCCESS	09	09	00.11	00.32	00.43	00.2	2 00.33	00.33

33	SUCCESS	11	11	00.11	00.1	00.21			0.33
24	SUCCESS	12	12	00.11	00.1	00.21			0.33
19	SUCCESS	14	14	00.11	00.1	00.21			0.33
08	SUCCESS	15	15	00.11	00.1	00.21			0.33
42	SUCCESS	16	16	00.11	00.1	00.21			0.33
11	SUCCESS	17	17	00.11	00.1	00.21	00 (0.11 0	0.33
20	SUCCESS	19	19	00.11	00.1	00.21	00 (0.11 0	0.33
27	SUCCESS	20	20	00.11	00.1	00.21	00 (00.11 0	0.33
16	SUCCESS	21	21	00.11	00.1	00.21	00 (0.11 0	0.33
40	SUCCESS	07	07	00.11	00.21	00.32	00.11	00.22	00.33
23	SUCCESS	09	09	00.11	00.21	00.32	00.11	00.22	00.33
47	SUCCESS	14	14	00.11	00.21	00.32	00.11	00.22	00.33
30	SUCCESS	15	15	00.11	00.21	00.32	00.11	00.22	00.33
45	SUCCESS	17	17	00.11	00.21	00.32	00.11	00.22	00.33
26	SUCCESS	19	19	00.11	00.21	00.32	00.11	00.22	00.33
39	SUCCESS	20	20	00.11	00.21	00.32	00.11	00.22	00.33
22	SUCCESS	21	21	00.11	00.21	00.32	00.11	00.22	00.33
35	SUCCESS	09	09	00.11	00.32	00.43	00.22	00.33	00.33
19	SUCCESS	14	14	00.11	00.32	00.43	00.22	00.33	00.33
32	SUCCESS	19	19	00.11	00.32	00.43	00.22	00.33	00.33
36	SUCCESS	09	09	00.11	00.43	00.54	00.33	00.44	00.33
34	SUCCESS	19	19	00.11	00.43	00.54	00.33	00.44	00.33
43	SUCCESS	09	09	00.11	00.54	00.65	00.44	00.55	00.33
48	SUCCESS	19	19	00.11	00.54	00.65	00.44	00.55	00.33
12	SUCCESS	05	05	184.97	00.1	185.07	00	184.97	554.91
41	SUCCESS	05	05	00.11	185.07	185.18	184.97	185.08	00.33
)5	SUCCESS	06	06	225.89	00.1	225.99	00	225.89	677.66
14	SUCCESS	06	06	00.11	225.99	226.1	225.89	226	00.33
17	SUCCESS	03	03	479.87	00.1	479.97	00	479.87	1439.6
)7	SUCCESS	02	02	513.88	00.1	513.98	00	513.88	1541.64
29	SUCCESS	02	02	322.75	513.98	836.73	513.88	836.63	
14	SUCCESS	04	04	1005.4	00.1	1005.5	00	1005.4	3016.19
37	SUCCESS	04	04	00.11	1005.5	1005.61	1005.4	1005.5	
snan II	sing RR: 1311.9	290535570512							

Total Cost : 8207.82600000001 Avg cost: 164.1565200000003

Avg Waiting Time: 38.664539999999995

4

org.cloudbus.cloudsim.examples.RoundRobinScheduler finished!

Makespan using RR: 1311.9290535570512

Total Completion Time: 4669.169 Avg Completion Time: 93.38338

Total Cost : 8207.82600000001 Avg cost: 164.1565200000003

Avg Waiting Time: 38.664539999999995

5.Shortest Job First :

rminated > S	JF_Scheduler (2) [Java	Application] C:\Us	sers\Shubham\Desl	ktop\cloud simsetup\	eclipse-java-2022-06-F	R-win32-x86_64\eclipse\plugins\org.eclips
27	SUCCESS	11	11	00.11	00.1	00.21
17	SUCCESS	13	13	00.11	00.1	00.21
37	SUCCESS	15	15	00.11	00.1	00.21
32	SUCCESS	16	16	00.11	00.1	00.21
33	SUCCESS	17	17	00.11	00.1	00.21
23	SUCCESS	19	19	00.11	00.1	00.21
28	SUCCESS	20	20	00.11	00.1	00.21
46	SUCCESS	05	05	00.11	00.21	00.32
44	SUCCESS	07	07	00.11	00.21	00.32
12	SUCCESS	08	08	00.11	00.21	00.32
08	SUCCESS	10	10	00.11	00.21	00.32
34	SUCCESS	11	11	00.11	00.21	00.32
24	SUCCESS	13	13	00.11	00.21	00.32
38	SUCCESS	15	15	00.11	00.21	00.32
43	SUCCESS	16	16	00.11	00.21	00.32
42	SUCCESS	17	17	00.11	00.21	00.32
36	SUCCESS	19	19	00.11	00.21	00.32
49	SUCCESS	20	20	00.11	00.21	00.32
21	SUCCESS	08	08	00.11	00.32	00.43
13	SUCCESS	10	10	00.11	00.32	00.43
47	SUCCESS	15	15	00.11	00.32	00.43
48	SUCCESS	16	16	00.11	00.32	00.43
45	SUCCESS	19	19	00.11	00.32	00.43
18	SUCCESS	10	10	00.11	00.43	00.54
41	SUCCESS	10	10	00.11	00.54	00.65
00	SUCCESS	06	06	1169.54	00.1	1169.64
11	SUCCESS	06	06	1095.08	1169.64	2264.72
02	SUCCESS	03	03	2920.54	00.1	2920.64
04	SUCCESS	03	03	2306.94	2920.64	5227.58
15	SUCCESS	03	03	1783.5	5227.58	7011.08
31	SUCCESS	03	03	00.11	7011.08	7011.19
39	SUCCESS	03	03	00.11	7011.19	7011.3

Makespan using SJF: 1490.0679364570929

Total Completion Time: 32621.78 Avg Completion Time: 652.4356 Total Cost: 27837.0660000006 Avg cost: 556.741320000001 Avg Waiting Time: 466.85515999999996 org.cloudbus.cloudsim.examples.SJF_Scheduler finished!

Datacenter_1 Broker_0 is S Simulation c	9 is shutti	ng down				
imulation c		wn				
Simulation co						
simulation c	Smpieled.					
	JTPUT ====					
loudlet ID	STATUS	Data center ID	VM ID	Time	Start Time	Finish Time
35	SUCCESS	05	05	00	00.1	00.1
03	SUCCESS	07	07	00	00.1	00.1
01	SUCCESS	08	08	00	00.1	00.1
05	SUCCESS	10	10	00	00.1	00.1
14	SUCCESS	11	11	00	00.1	00.1
19	SUCCESS	12	12	00	00.1	00.1
16	SUCCESS	13	13	00	00.1	00.1
29	SUCCESS	15	15	00	00.1	00.1
20	SUCCESS	16	16	00	00.1	00.1
30	SUCCESS	17	17	00	00.1	00.1
25	SUCCESS	18	18	00	00.1	00.1
22	SUCCESS	19	19	00	00.1	00.1
09	SUCCESS	20	20	00	00.1	00.1
26	SUCCESS	21	21	00	00.1	00.1
40	SUCCESS	05	05	00.11	00.1	00.21
07	SUCCESS	07	07	00.11	00.1	00.21
10	SUCCESS	08	08	00.11	00.1	00.21
06	SUCCESS	10	10	00.11	00.1	00.21
27	SUCCESS	11	11	00.11	00.1	00.21
17	SUCCESS	13	13	00.11	00.1	00.21
37	SUCCESS	15	15	00.11	00.1	00.21
32	SUCCESS	16	16	00.11	00.1	00.21
33	SUCCESS	17	17	00.11	00.1	00.21
23	SUCCESS	19	19	00.11	00.1	00.21
28	SUCCESS	20	20	00.11	00.1	00.21
46	SUCCESS	05	05	00.11	00.21	00.32
44	SUCCESS	07	07	00.11	00.21	00.32
12	SUCCESS	08	08	00.11	00.21	00.32
08	SUCCESS	10	10	00.11	00.21	00.32

Makespan using SJF: 1490.0679364570929 Total Completion Time: 32621.78 Avg Completion Time: 652.4356 Total Cost : 27837.06600000006 Avg cost: 556.741320000001 Avg Waiting Time: 466.85515999999996

6. Selective Algortihm

0.1: Broker: Sending cloudlet 8 to VM #3 oo it miloc xoo_or(cenpsetpidgilistorgieenps 0.1: Broker: Sending cloudlet 9 to VM #3 54.5266666666666664: Broker: Cloudlet 6 received 63.41555555555556: Broker: Cloudlet 4 received 71.19333333333333: Broker: Cloudlet 7 received 77.86: Broker: Cloudlet 5 received 83.41555555555556: Broker: Cloudlet 9 received 87.86: Broker: Cloudlet 0 received 91.19333333333333: Broker: Cloudlet 8 received 93.41555555555556: Broker: Cloudlet 1 received 95.63777777777779: Broker: Cloudlet 2 received 95.6377777777779: Broker: All Cloudlets executed. Finishing... 95.63777777777779: Broker: Destroying VM #0 95.63777777777779: Broker: Destroying VM #1 95.63777777777779: Broker: Destroying VM #2 95.63777777777779: Broker: Destroying VM #3 Broker is shutting down... Simulation: No more future events CloudInformationService: Notify all CloudSim entities for shutting down. Datacenter 0 is shutting down... Broker is shutting down... Simulation completed. Simulation completed. ======= OUTPUT ========
 DUTPUT
 STATUS
 Data center ID
 VM ID
 Time
 Start Time
 Finish Time
 Algo

 3
 SUCCESS
 2
 3
 44.44
 0.1
 44.54
 Max-min

 6
 SUCCESS
 2
 3
 54.43
 0.1
 54.53
 Max-min

 4
 SUCCESS
 2
 3
 63.32
 0.1
 63.42
 Max-min

 7
 SUCCESS
 2
 3
 63.32
 0.1
 63.42
 Max-min

 5
 SUCCESS
 2
 3
 71.09
 0.1
 71.19
 Max-min

 9
 SUCCESS
 2
 3
 83.32
 0.1
 83.42
 Min-min

 0
 SUCCESS
 2
 3
 87.76
 0.1
 87.86
 Max-min

 8
 SUCCESS
 2
 3
 91.09
 0.1
 91.19
 Max-min

 1
 SUCCESS
 2
 3
 93.32
 0.1
 93.42
 Max-min

 2
 SUCCESS
 2

Selective Algorithm finished!

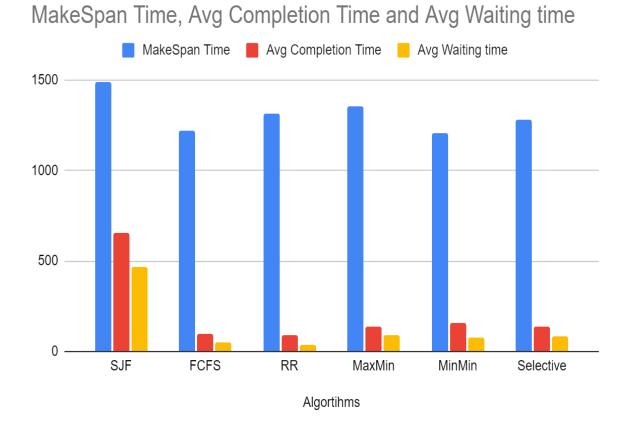
4.2 Result and Comparison:

There is the comparison between Various Task scheduling algorithms.

Algortihms	MakeSpan Time	AvgCompletion Time	Avg Waiting time
SJF	1490.06	652.43	466.85
FCFS	1221.22	98.95	48.03
RR	1311.92	93.38	38.66
MaxMin	1354.52	139.57	93.62
MinMin	1208.41	155.88	80.07
Selective	1280.54	140.45	81.34

The results above show that considering MakeSpan time MinMin Algorithms is efficient followed by FCFS algorithms.

- Taking Avg completion RR algorithms shows the best result followed by FCFS.
- Taking account Avg Waiting time RR algorithm is the best followed by FCFS.





Scheduling Sim	ulator						
Input		Choose File	lo file chosen	Submit	Console Output		
Task Execution Time 1	Period	Deadline	Arrival Time	Add Row			
MxMn RR MnMn Edf	FCFS SJF						
Max-Min					Bar Graph	Line Graph	

Fig 4.6 Scheduling Simulator

	4	· · · · · · · · · · · · · · · · · · ·
MxMn RR MnMn Edf FCFS SJF		
Max-Min	Bar Graph Line Graph	C D
Round Robin	Bar Graph Line Graph	
Min Min	Bar Graph Line Graph	C 2
		Ţ

CONCLUSION

In the project, we have introduced the main concepts of schedulings and resources management in the cloud computing. We have presented a comparative analysis of schedulings algorithm by considering the factors waiting time , execution cost ,most important makespan time (execution time). The problem of task scheduling in the Cloud Computing environment is concerned .Maxmin,Minmin,Round Robin are famous algorithms for scheduling tasks on cloud computing environments. Minmin,Maxmin algorithms are scheduled to reduce the make-span time taken , average-waiting time and execution cost. The comparison result indicate that RR,Maxmin,Minmin showed the better results with less overhead.

Future Work

- Future work should examine the impact of additional factors like virtual machines (VMs), data centers, memory, network bandwidth, and storages in cloud environment and reflect those factors in actual physical environments.
- In future works we have planned to improve upon the current algorithms and implement best algorithm in order optimize makespan time, total execution cost.
- We plan to create a webapp to better implementation of your project.
- The project intend to develop the suggested system by extending the practical Cloud platform, such as OpenStack, in order to assess it in a genuine Cloud architecture.
- This work has significant social significance because it lowers modern IT infrastructures' carbon dioxide footprint and energy consumptions in addition to lowering infrastructure and ongoing operating costs.

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```
package org.cloudbus.cloudsim.examples;
@import org.cloudbus.cloudsim.*;[]
 public class SJF Scheduler {
     private static List<Cloudlet> cloudletList;
     private static List<Vm> vmList;
     private static Datacenter[] datacenter;
     private static double[][] commMatrix;
     private static double[][] execMatrix;
    private static List<Vm> createVM(int userId, int vms) {
Θ
         //Creates a container to store VMs. This list is passed to the broker later
        LinkedList<Vm> list = new LinkedList<Vm>();
         //VM Parameters
         long size = 10000; //image size (MB)
         int ram = 512; //vm memory (MB)
         int mips = 250;
        long bw = 1000;
         int pesNumber = 1; //number of cpus
        String vmm = "Xen"; //VMM name
        //create VMs
        Vm[] vm = new Vm[vms];
         for (int i = 0; i < vms; i++) {</pre>
            vm[i] = new Vm(datacenter[i].getId(), userId, mips, pesNumber, ram, bw, size, vmm, new CloudletSchedulerSpaceShared());
             list.add(vm[i]);
         }
         return list;
     }
     private static List<Cloudlet> createCloudlet(int userId, int cloudlets, int idShift) {
Θ
         // Creates a container to store Cloudlets
         LinkedList<Cloudlet> list = new LinkedList<Cloudlet>();
        //cloudlet parameters
        long fileSize = 300;
        long outputSize = 300;
         int pesNumber = 1;
        UtilizationModel utilizationModel = new UtilizationModelFull();
         Cloudlet[] cloudlet = new Cloudlet[cloudlets];
                                                                                                  Smart Incert 38 · 52 · 1415
                                                                                     Writable
```

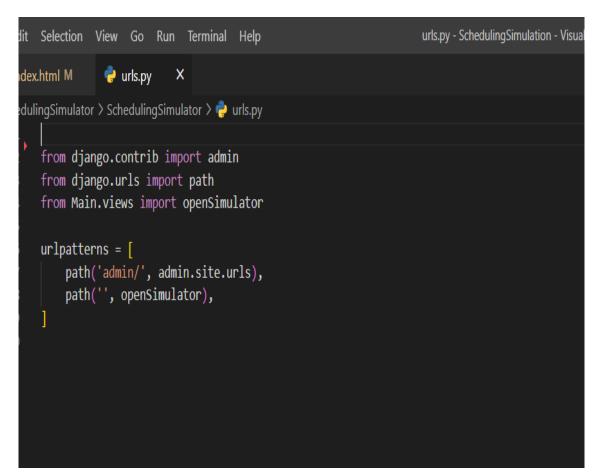
```
package org.cloudbus.cloudsim.examples;
mport org.cloudbus.cloudsim.*;[]
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     private static List<Vm> vmList;
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     private static double[][] commMatrix;
     private static double[][] execMatrix;
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             vm[i] = new Vm(datacenter[i].getId(), userId, mips, pesNumber, ram, bw, size, vmm, new CloudletSchedulerSpaceShared());
             list.add(vm[i]);
         }
         return list;
     }
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         long fileSize = 300;
         long outputSize = 300;
         int pesNumber = 1;
         UtilizationModel utilizationModel = new UtilizationModelFull();
         Cloudlet[] cloudlet = new Cloudlet[cloudlets];
 4
                                                                                                                      28 - 52 - 1415
                                                                                      Writable
                                                                                                       Smart Incert
```

4	from django.shortcuts import render
5 6	from django.views.decorators.csrf import csrf_protect
7	from .forms import UserCreationForm
8	from Main import schedulingAlgorithm
9	
10	@csrf_protect
11	def openSimulator(request) :
12	# for printing the console into a variable
13	old_stdout = sys.stdout
14	<pre>new_stdout = io.StringIO()</pre>
15	sys.stdout = new_stdout
16	
17	<pre>if request.method == 'POST':</pre>
18	form = UserCreationForm()
19	<pre>final_array = request.POST.get('final_array')</pre>
20	<pre>interval = request.POST.get('interval')</pre>
21	<pre>noOfTasks = request.POST.get('noOfTasks')</pre>
22	#form['interval'].initial = interval
23	
24	# splitting the input into desired format
25	algorithmInputArray = arraySplit(final_array)
26	# calculating rate monotonic scheduling algorithm
27	rmsArray, edfArray, fcfsArray, sjfArray = schedulingAlgorithm.schedule(algorithmInputArray)
28	# formating the input array to send it back to screen
29	algorithmInputArray = returnInputArray(algorithmInputArray)

writing standard output to variable
<pre>outputConsole = new_stdout.getvalue()</pre>
<pre>sys.stdout = old_stdout # setting the standard output back to console</pre>
return render(request, 'index.html',{'sjfArray': sjfArray, 'fcfsArray': fcfsArray, 'edfArray': edfArray, 'rmsArray, 'noOfTasks
else:
form = UserCreationForm()
<pre>outputConsole = new_stdout.getvalue()</pre>
<pre>sys.stdout = old_stdout # setting the standard output back to console</pre>
return render(request, 'index.html', {'sjfArray': [], 'fcfsArray': [], 'edfArray': [], 'rmsArray': [], 'noOfTasks': [], 'algorithmInputArr
<pre>def arraySplit(algorithmInputs): algorithmInputs = algorithmInputs[:-1] algorithmInputs = algorithmInputs.split(" ") algorithmInputArray = [] for algorithmInput in algorithmInputs: algorithmInput = algorithmInput[:-1] algorithmInputArray.append(list(map(int, algorithmInput.split(",", -1)))) return algorithmInputArray</pre>
def returnInputArray(algorithmInputArray): algorithmInputArray.sort(key=lambda x: x[1]) for inputArray in algorithmInputArray: inputArray[1:] return algorithmInputArray

```
MaxminBroke... D Maxminjava D Selective.java D MinMinjava D RoundRobinSc... D FCFS_Schedul...
package Utils;

public class Constants {
    public static final int NO_OF_TASKS = 50; // number of Cloudlets;
    public static final int NO_OF_DATA_CENTERS = 20; // number of Datacenters;
    public static final int POPULATION_SIZE = 25; // Number of Particles.
    }
```



Major Report

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3	slidepla	yer.com			<1,	
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